
BITS AND FLOPS IN NON-ERGODIC MIMO:

CAN (A REASONABLE NUMBER OF) FLOPS PROVABLY
OFFER ERGODIC-LIKE BEHAVIOR

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Performance and complexity in MIMO communications

- Setting of interest: general *outage limited* MIMO communications

$$\mathbf{y} = \mathbf{H}\mathbf{x} + \mathbf{w}$$

- ★ MIMO, MIMO-OFDM, MIMO-MAC, MIMO-ARQ, COOPERATIVE, HYBRID...
- ★ Rx knows \mathbf{H} , Tx does not

Examples:

- Communication of CSIT over feedback link (even with reciprocity in multiuser case)
- After interference cancellation

Performance and complexity in MIMO communications₁

(SNR ρ , rate R , reliability P_{err} , complexity C)

Rate and reliability measure and exponent

$$P_{\text{err}} = P(\hat{\mathbf{x}} \neq \mathbf{x}_{\text{tx}})$$

$$R = \frac{1}{T} \log |\text{Code}|, \quad |\text{Code}| = 2^{RT}$$

HIGH-SNR P_{ERR} BEHAVIOR: EXPONENT OVER ρ

$$d(r) := - \lim_{\rho \rightarrow \infty} \frac{\log P_{\text{err}}}{\log \rho}, \quad P_{\text{err}} \doteq \rho^{-d(r)} \quad r = \frac{R}{\log \rho}$$

Computational complexity measure

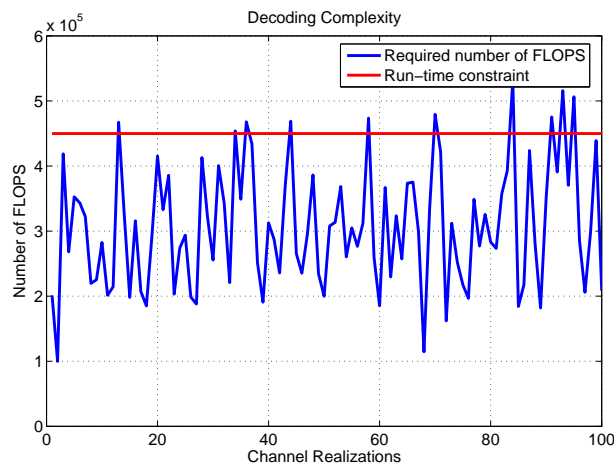
$$N_{\max}$$

MAXIMUM ALLOWABLE COMPUTATIONAL RESOURCES (PER T CHANNEL USES)

- chip size, number of flops (after that effort must terminate), etc.

Fluctuating complexity introduces a tradeoff

- Keep in mind: Generally complexity fluctuates with channel
- Generally $P_{\text{err}} \uparrow$ as $N_{\text{max}} \downarrow$



Instantaneous algorithmic complexity fluctuations

SMALL EXAMPLE

- Can you achieve $(P_{\text{err}}, R, \rho)$ with $N_{\text{max}} = 2000$ flops?
 - ★ *No! Too common early-terminations for search based decoders ($N(H)$ varies) - or too weak linear receivers*
- Can you do it with $N_{\text{max}} = 100000$ flops?
 - ★ *No, but we are getting there.*
- How about with 132957 flops?
 - ★ *Yes!*
- How about with 132956 flops?
 - ★ *No!*
- OK, for $(P_{\text{err}}, R, \rho)$ you need $N_{\text{max}} = 132957$ flops. Else $(P_{\text{err}}, R, \rho)$ is not achievable.

Complexity exponent

$$c(r) := \lim_{\rho \rightarrow \infty} \frac{\log N_{\max}}{\log \rho},$$

$$N_{\max} \doteq \rho^{c(r)} = 2^{R \frac{c(r)}{r}} \leq \rho^{rT} = |\mathcal{X}|$$

$c(r) > 0 \implies N_{\max}$ exponential in R (and often in RT)

Meaningful matching of error and complexity exponents

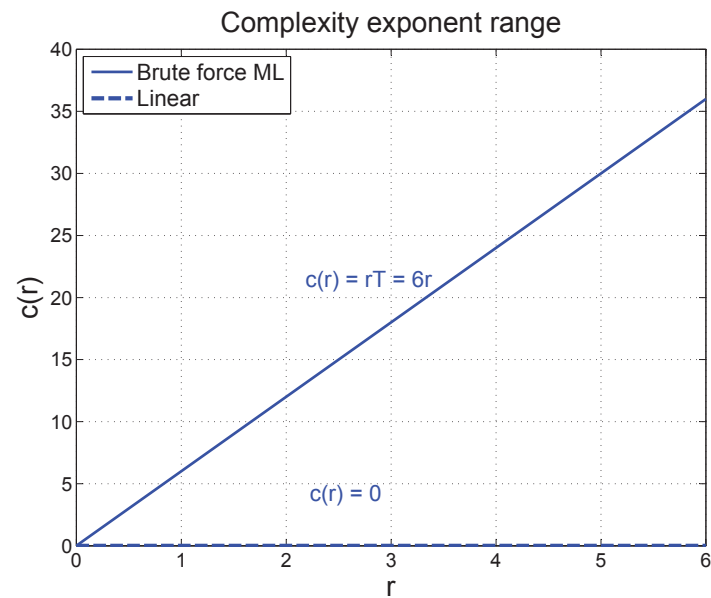
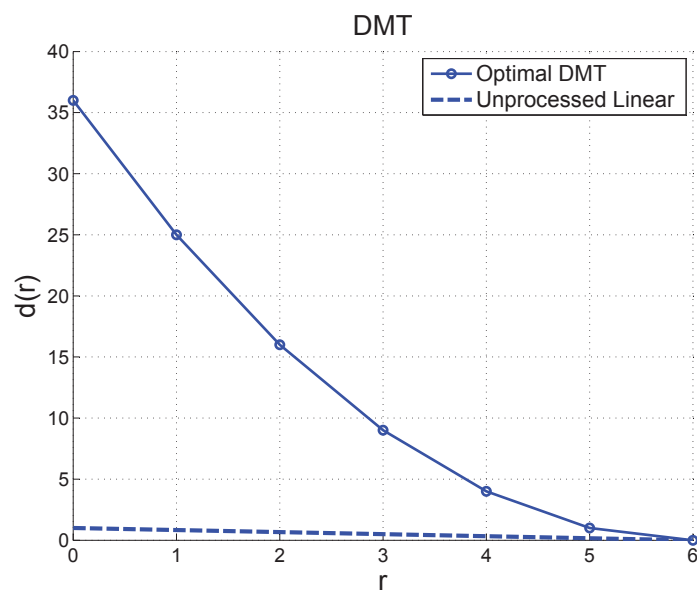
$$c(r) := \lim_{\rho \rightarrow \infty} \frac{\log N_{\max}}{\log \rho}, \quad d(r) := - \lim_{\rho \rightarrow \infty} \frac{\log P_{\text{err}}}{\log \rho}$$

- Reliability and complexity naturally polynomial in ρ

$$N_{\max} : 1 \rightarrow K \cdot |\text{Code}| \approx 2^{RT} \approx \rho^{rT}, \quad P_{\text{err}} : \rho^0 \rightarrow \rho^{-d_{\text{opt}}(r)}$$

Practical ramifications of both exponents

- Performance: From highly unreliable to near-ergodic reliability
- Complexity: From easy to impossible
 - ★ $c(r) = 0$: linear - very fast
 - ★ $c(r) = rT \longrightarrow N_{\max} = 2^{rT} \rightarrow \rho^{36}$



For now focus on search-based lattice decoders

- For now we neglect linear receivers
 - ★ Without lattice reduction they are extremely suboptimal
 - ★ May have unbounded gap to optimal solutions even with LR
 - ★ LR problematic in ubiquitous scenarios (inner-outer codes)
- We instead focus on search based (ML and lattice decoding)
 - ★ aka: ML sphere decoding, and lattice sphere decoding
- We also focus on linear lattice code designs

Search-based decoders

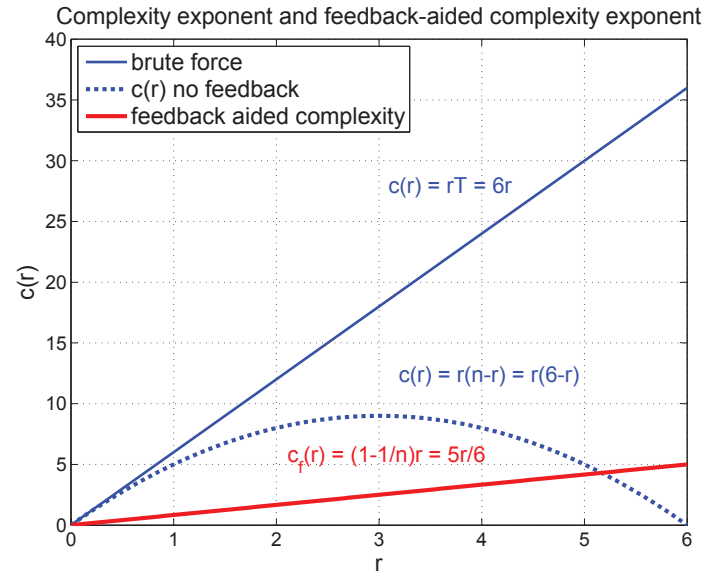
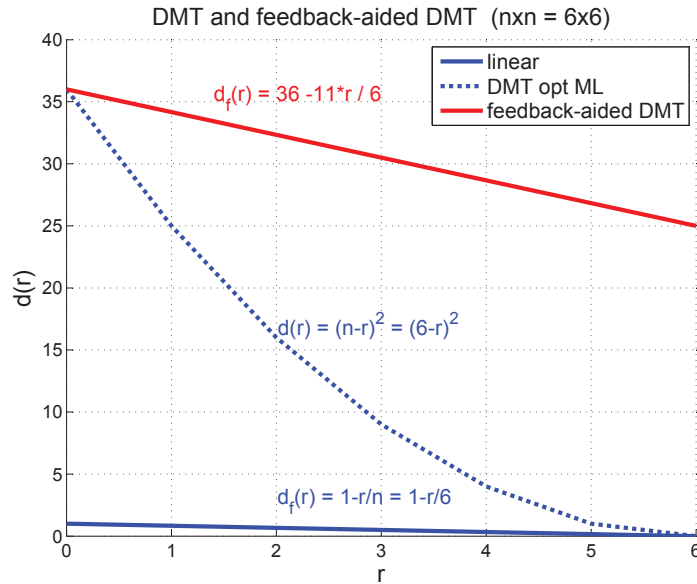
SEARCH-BASED DECODERS

ML-BASED SPHERE DECODING

(REGULARIZED) LATTICE-BASED SPHERE DECODING

NO LATTICE REDUCTION

Preview: $d(r), c(r)$ for ML and lattice decoders



- Answer will lie somewhere in the middle

Search based decoders

$$\mathbf{y} = \mathbf{H}\mathbf{x} + \mathbf{w} = \mathbf{Q}\mathbf{R}\mathbf{x} + \mathbf{w}$$

ML: Solve

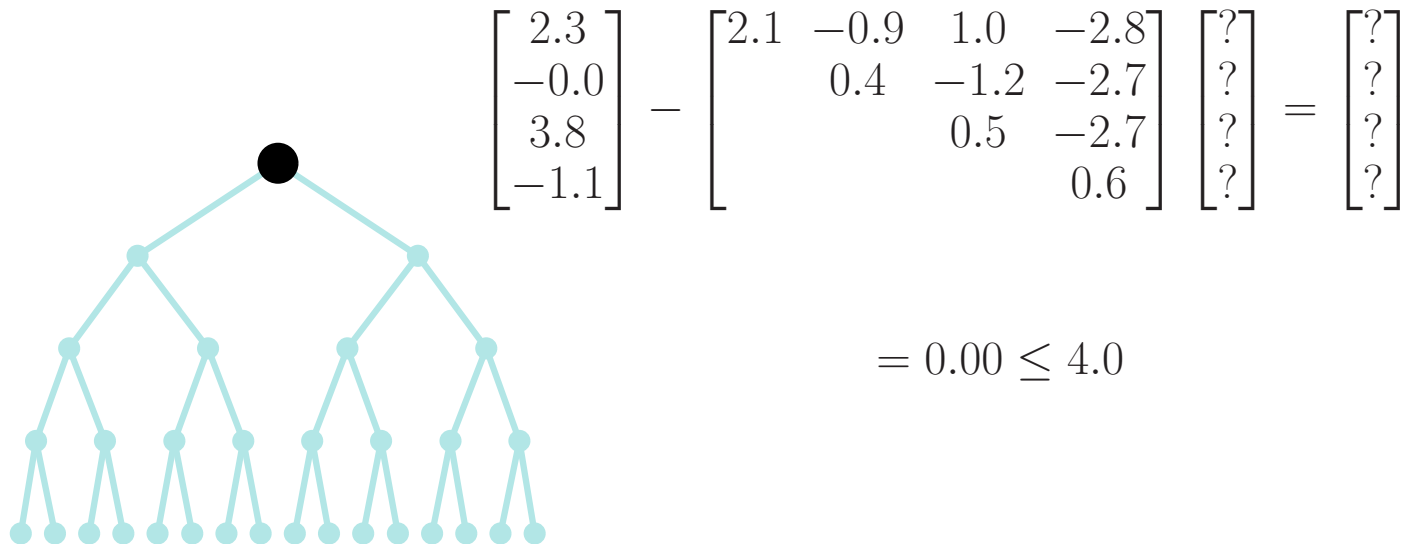
$$\min_{\hat{\mathbf{x}} \in \text{Code}} \|\tilde{\mathbf{y}} - \mathbf{R}\hat{\mathbf{x}}\|^2$$

(Regularized) Lattice decoding: Solve

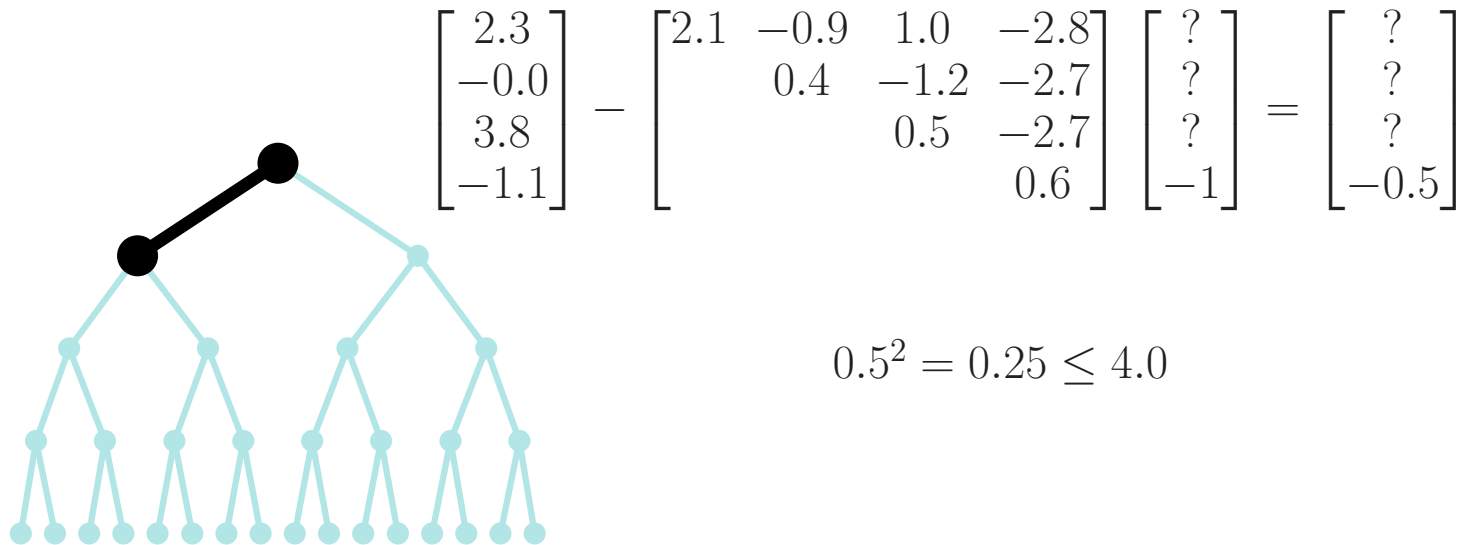
$$\min_{\hat{\mathbf{x}} \in \text{Lattice}} \|\tilde{\mathbf{y}} - \mathbf{R}\hat{\mathbf{x}}\|^2 + \alpha \|\hat{\mathbf{x}}\|_T^2$$

by searching over $\|\tilde{\mathbf{y}} - \mathbf{R}\hat{\mathbf{x}}\|^2 \leq (\text{radius})^2$

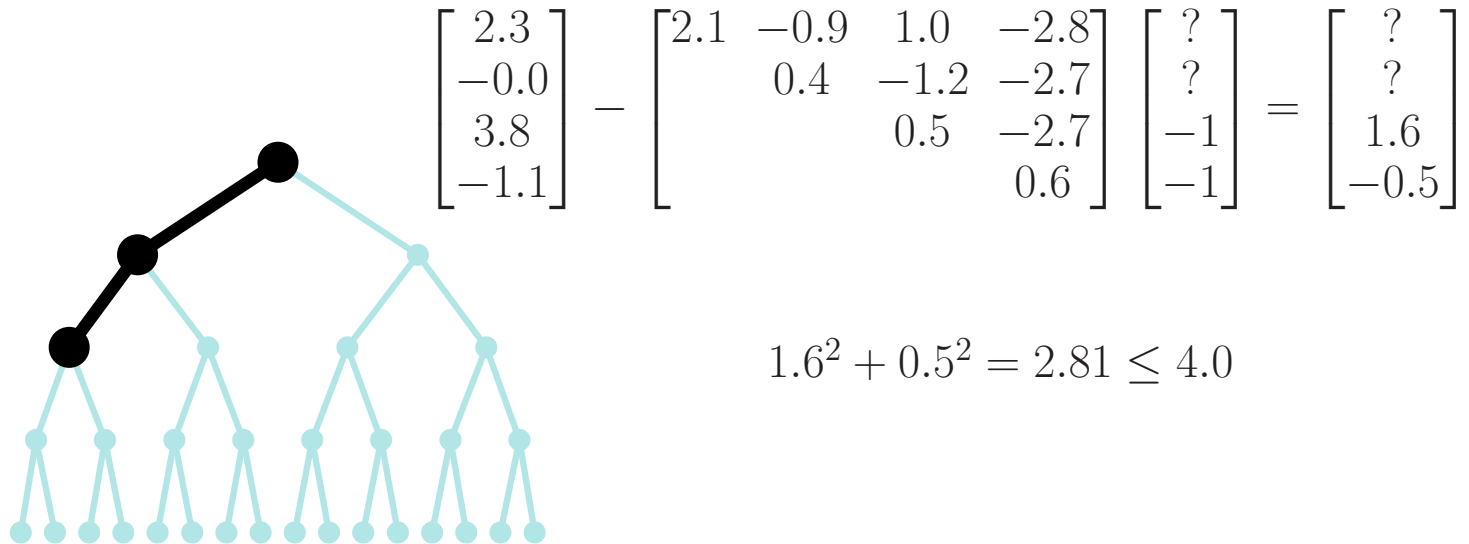
Sphere decoding example, numerics



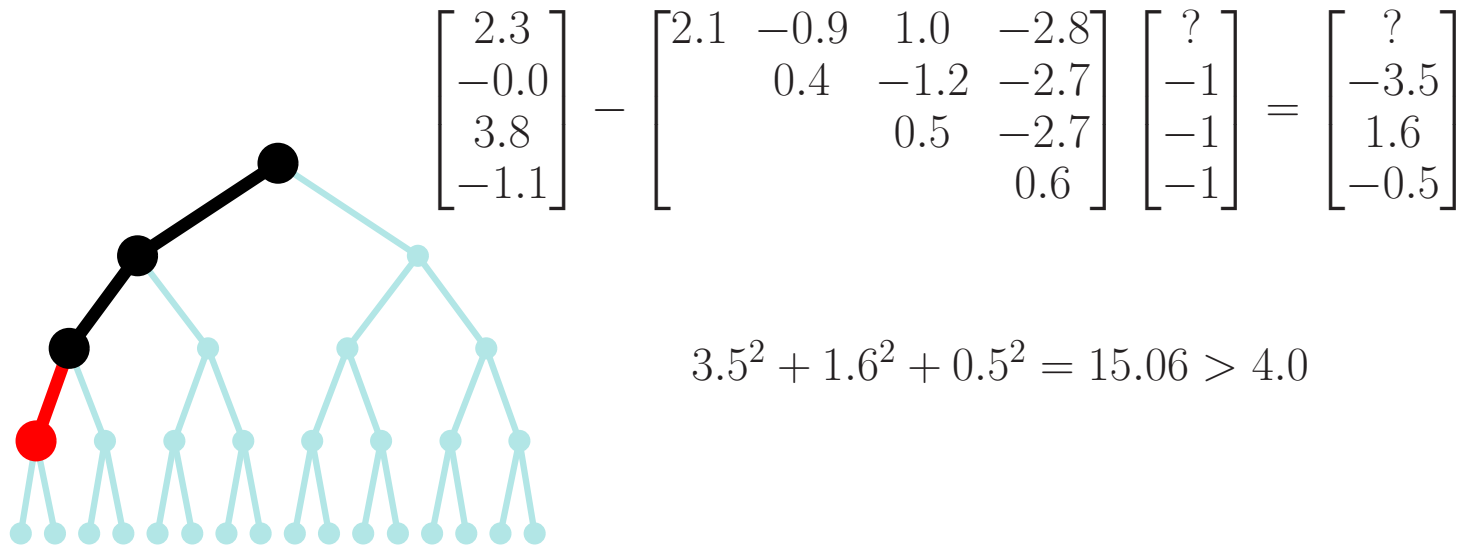
Sphere decoding example, numerics



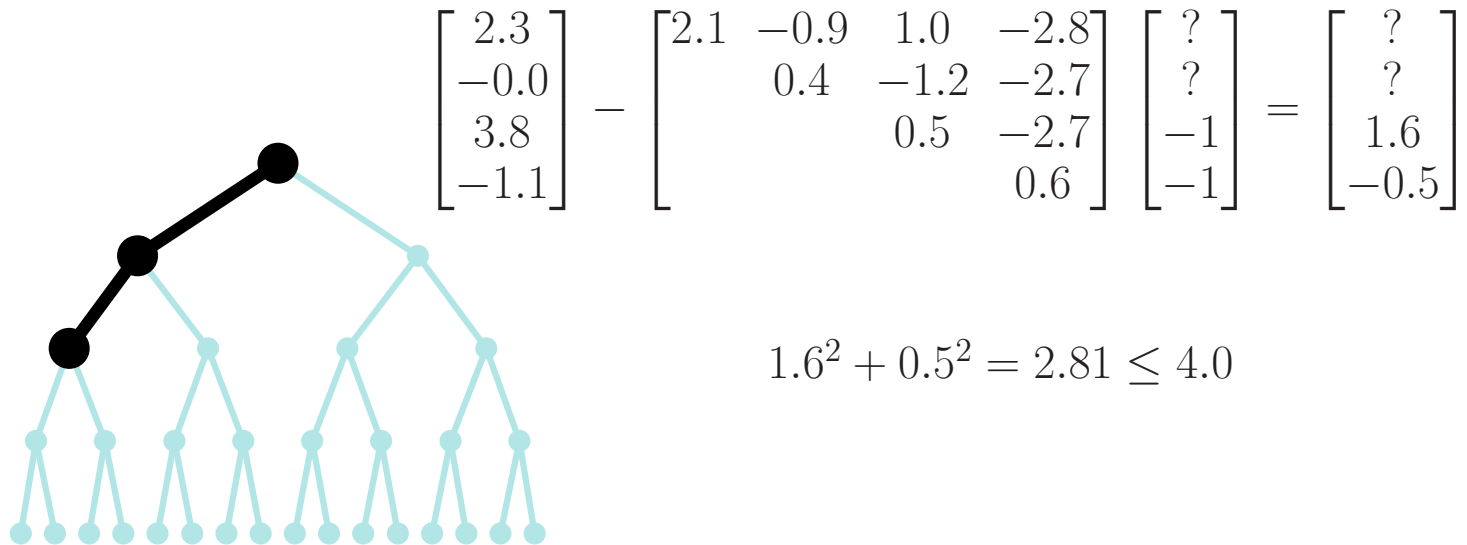
Sphere decoding example, numerics



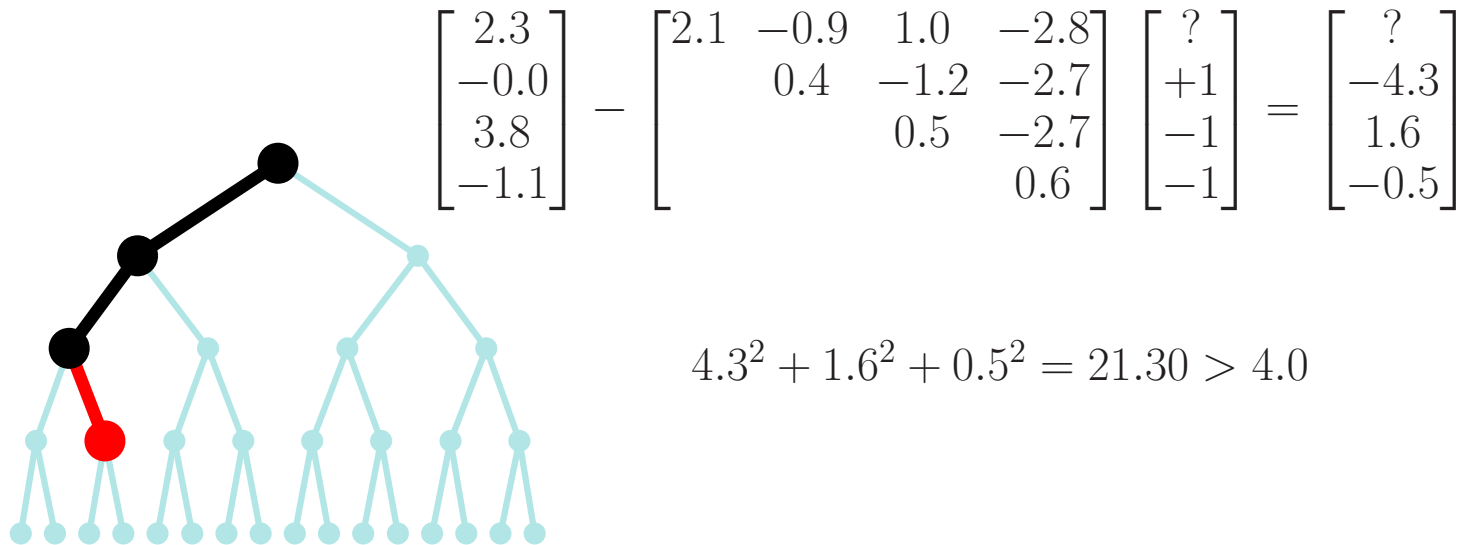
Sphere decoding example, numerics



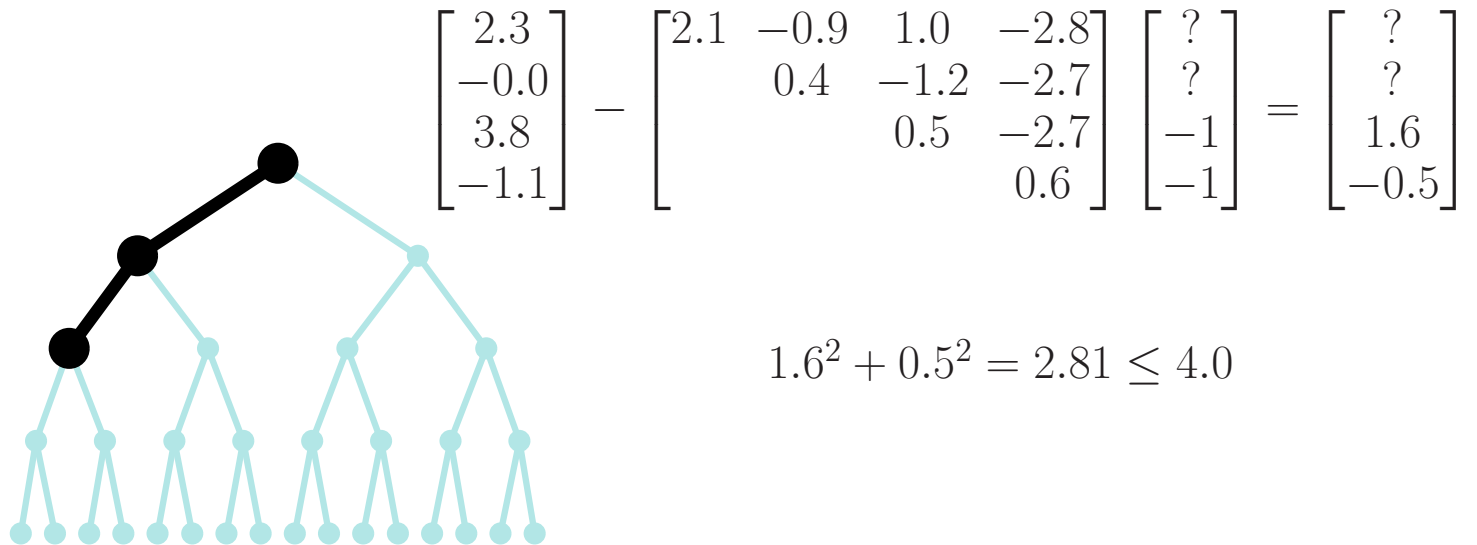
Sphere decoding example, numerics



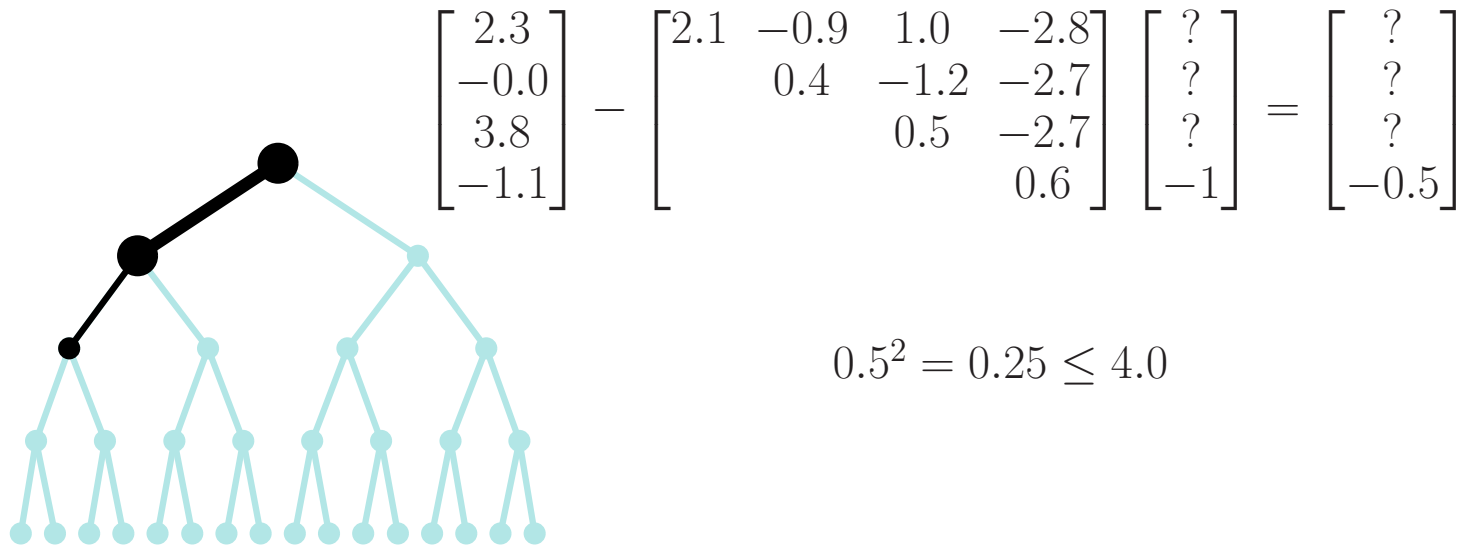
Sphere decoding example, numerics



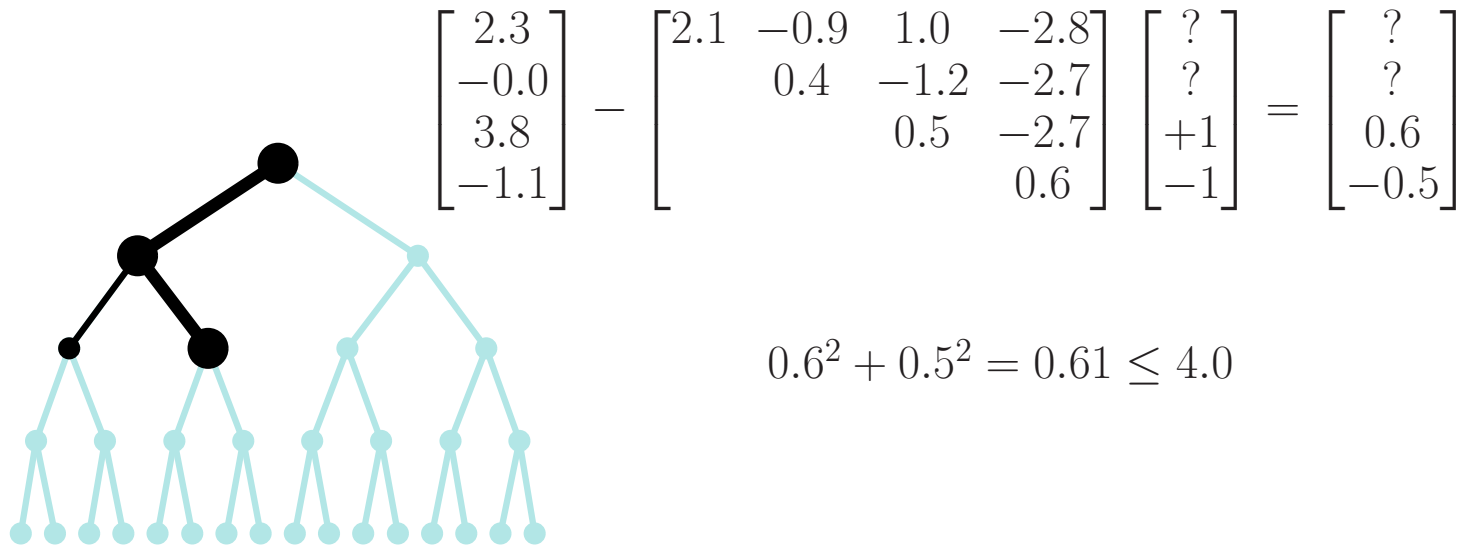
Sphere decoding example, numerics



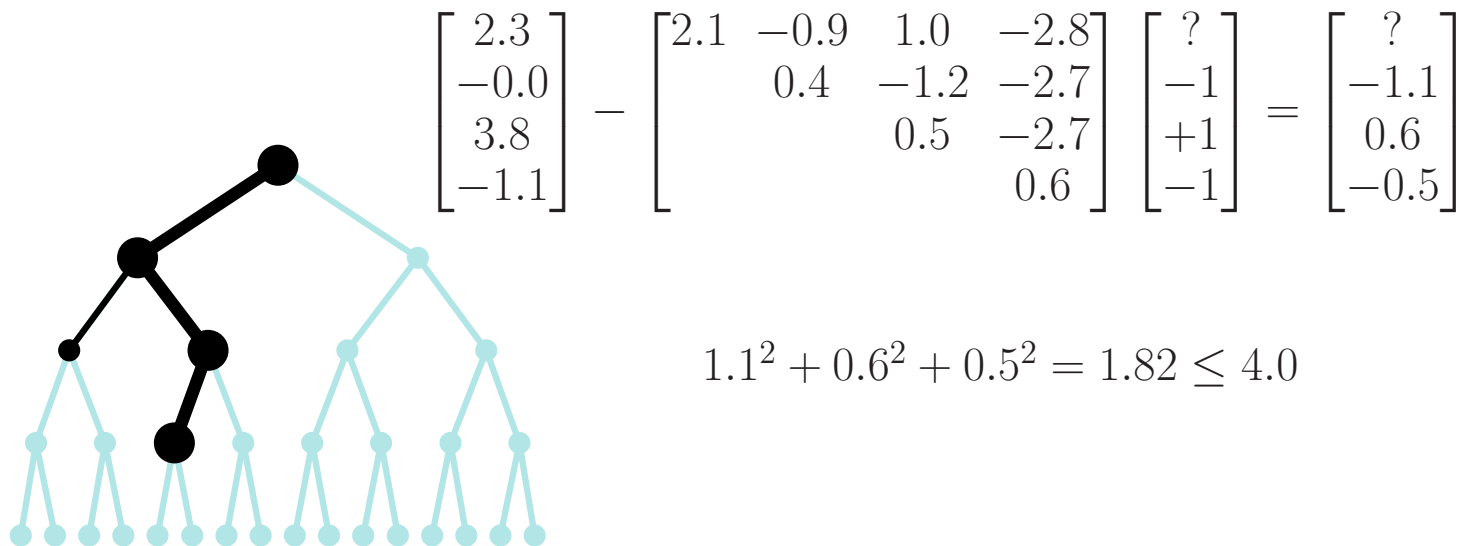
Sphere decoding example, numerics



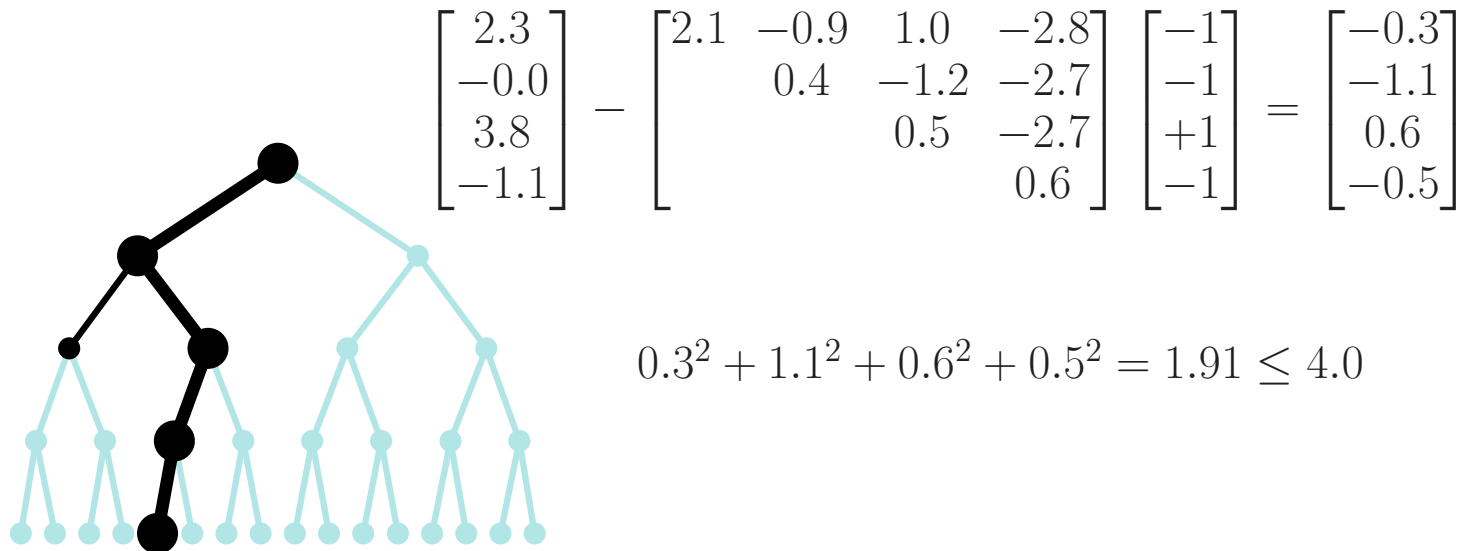
Sphere decoding example, numerics



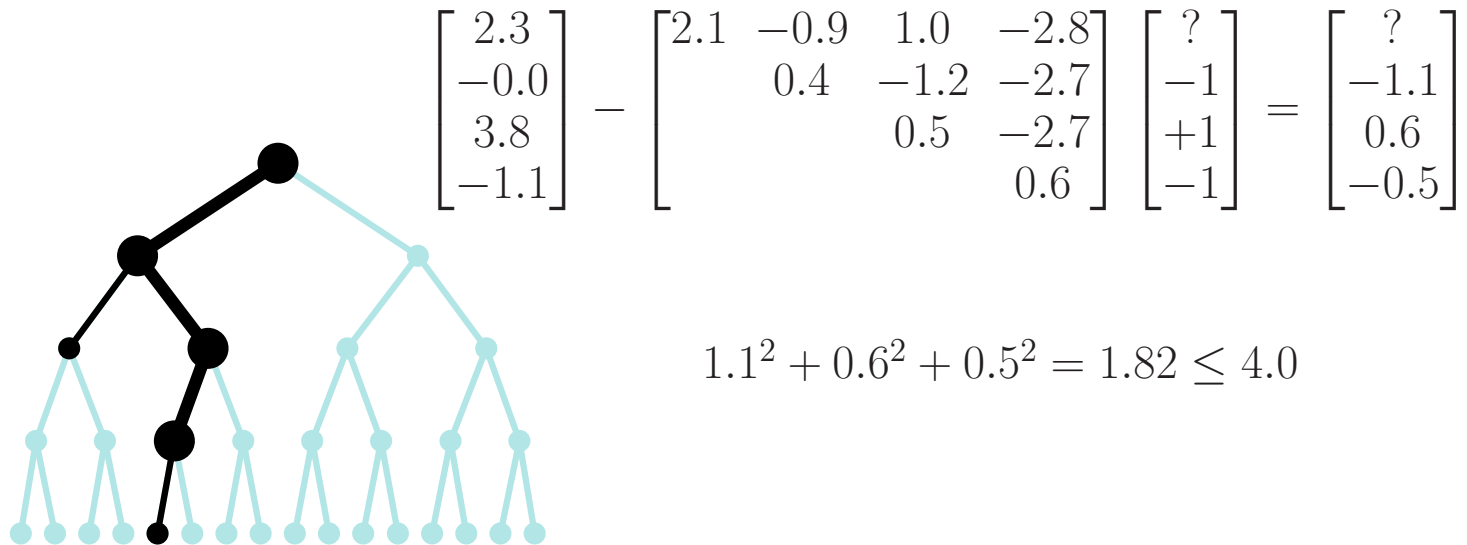
Sphere decoding example, numerics



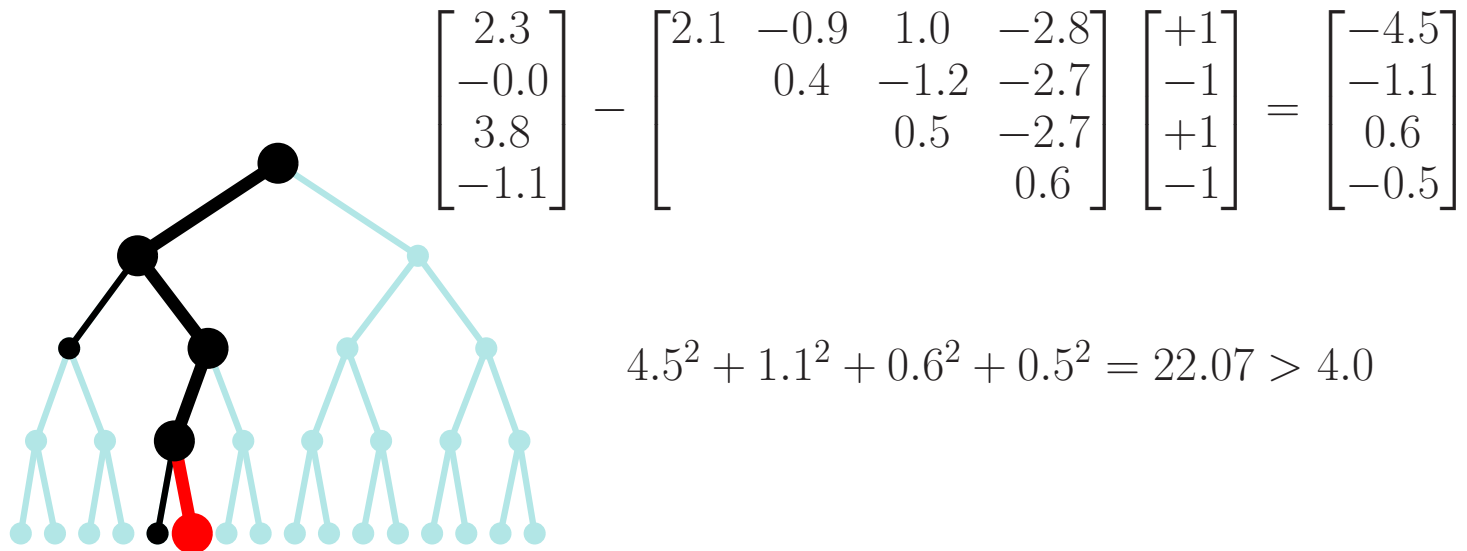
Sphere decoding example, numerics



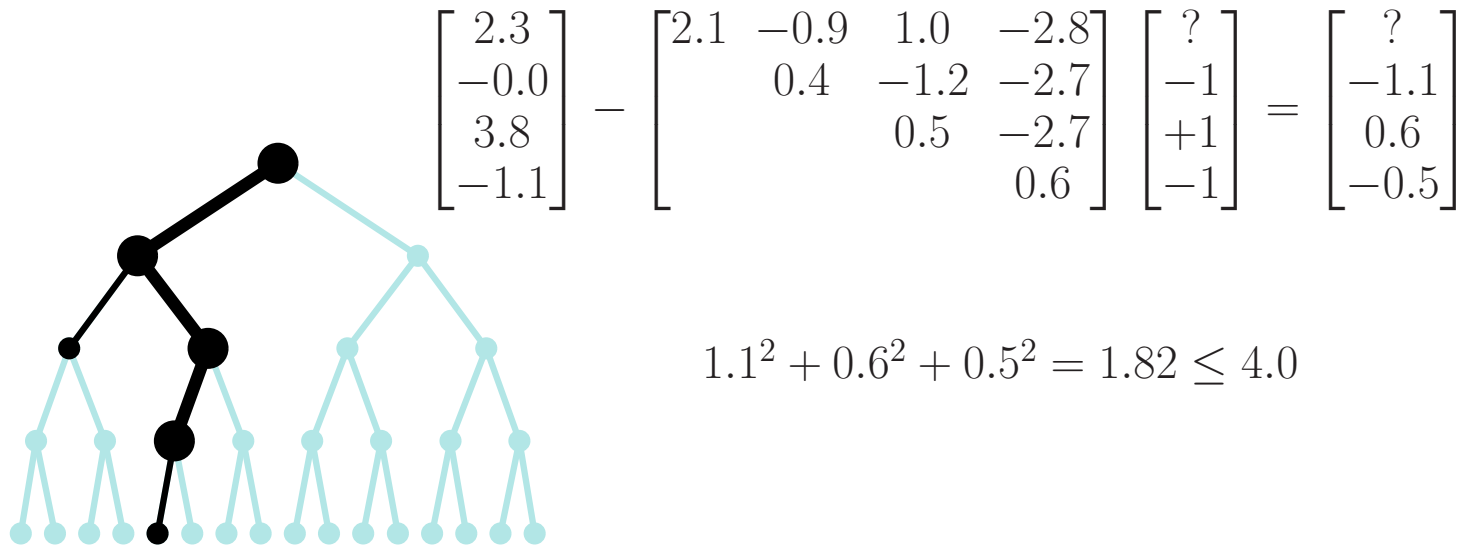
Sphere decoding example, numerics



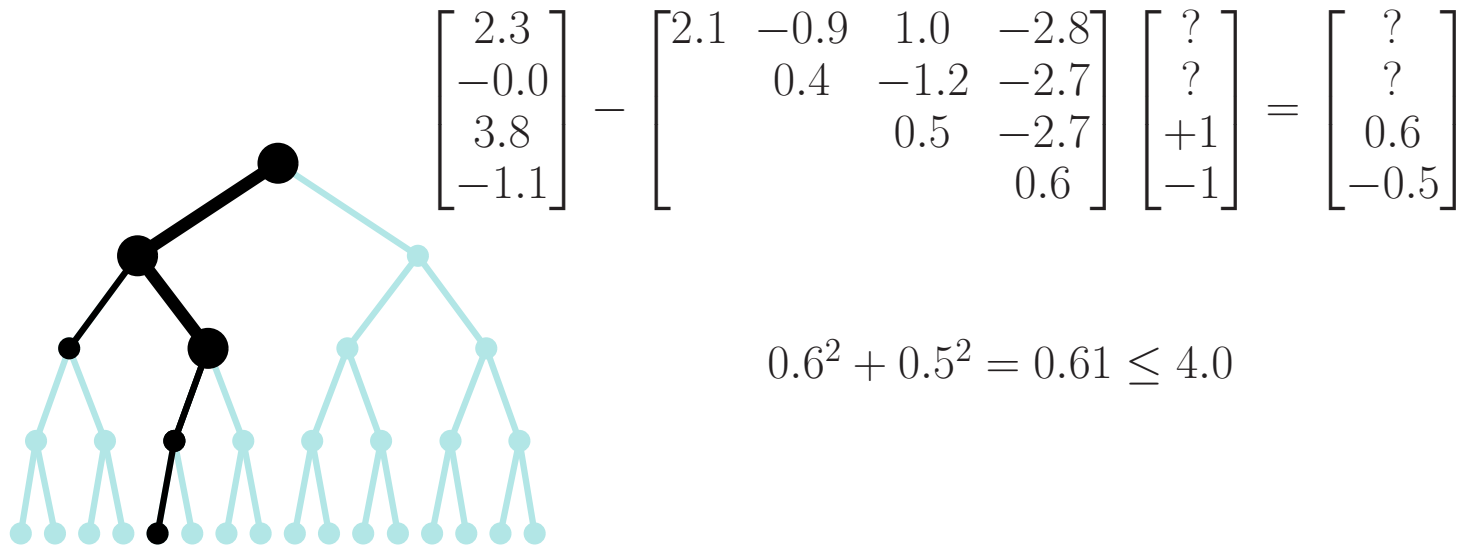
Sphere decoding example, numerics



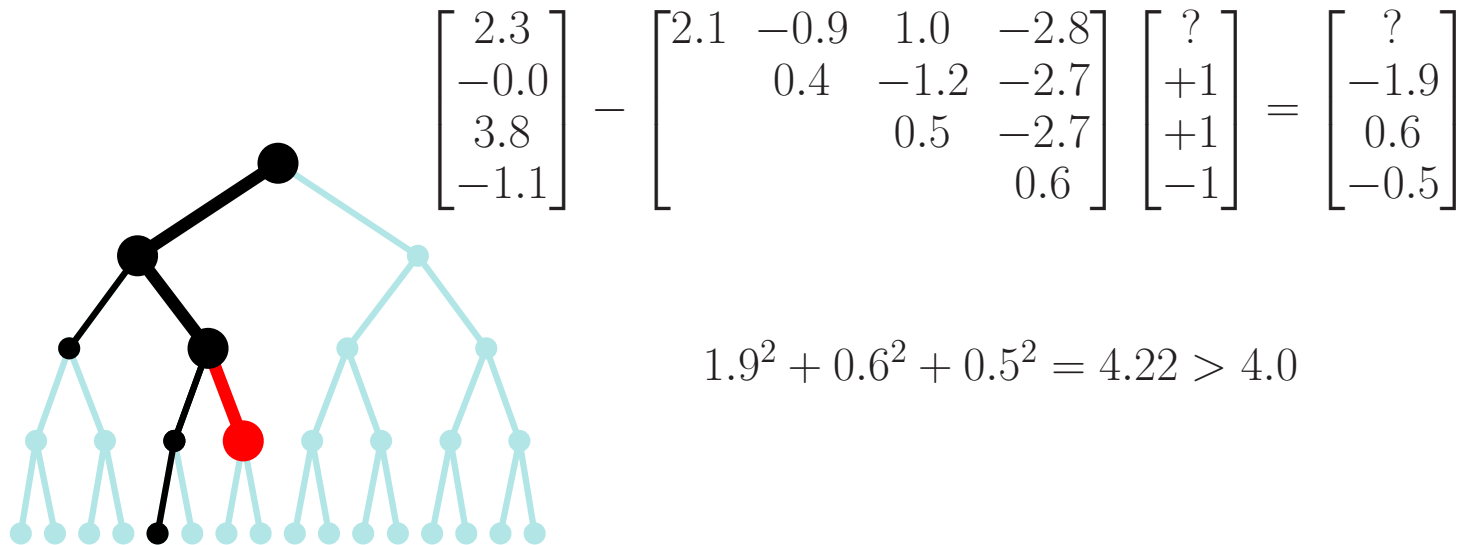
Sphere decoding example, numerics



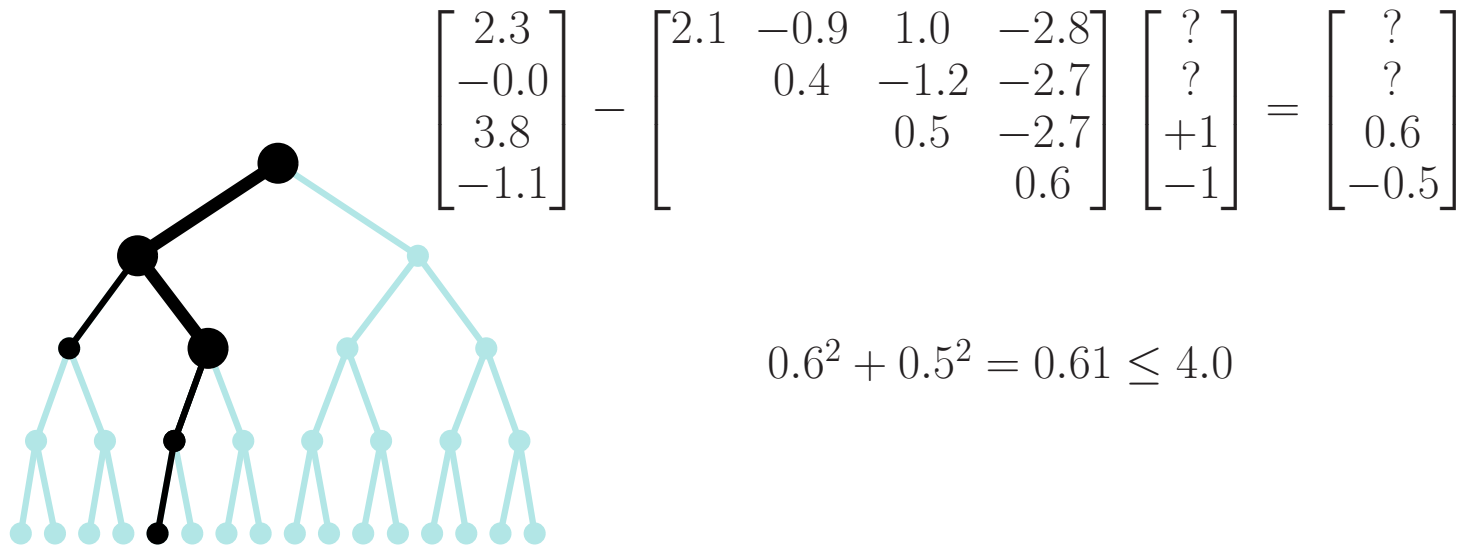
Sphere decoding example, numerics



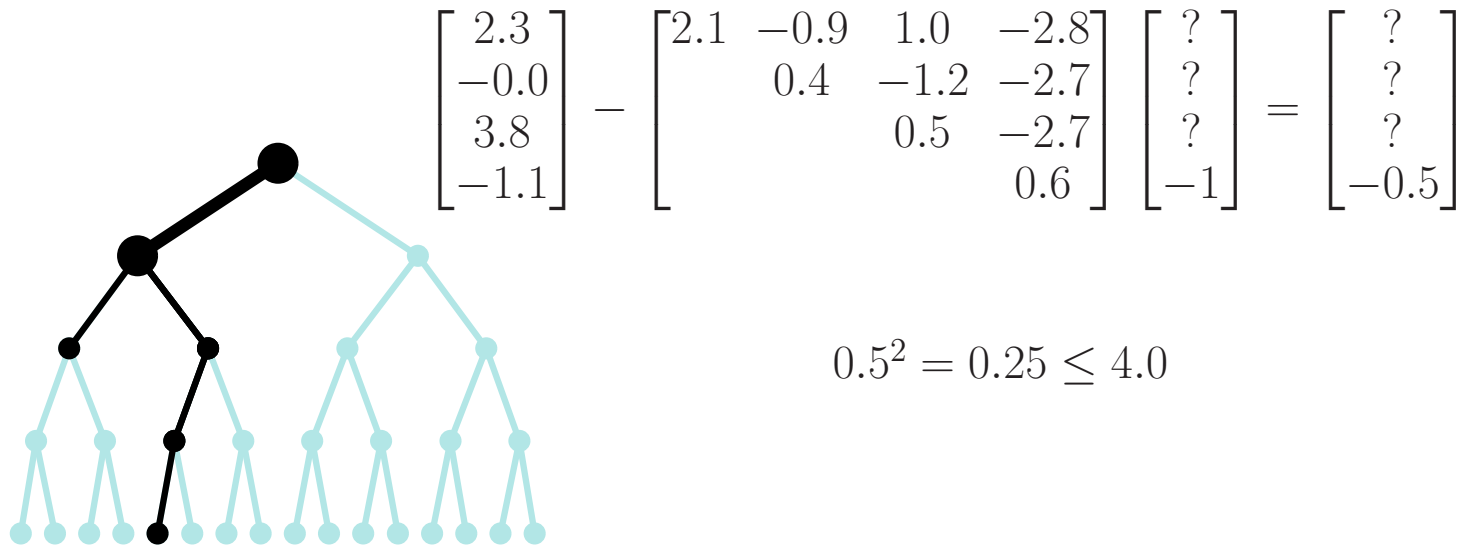
Sphere decoding example, numerics



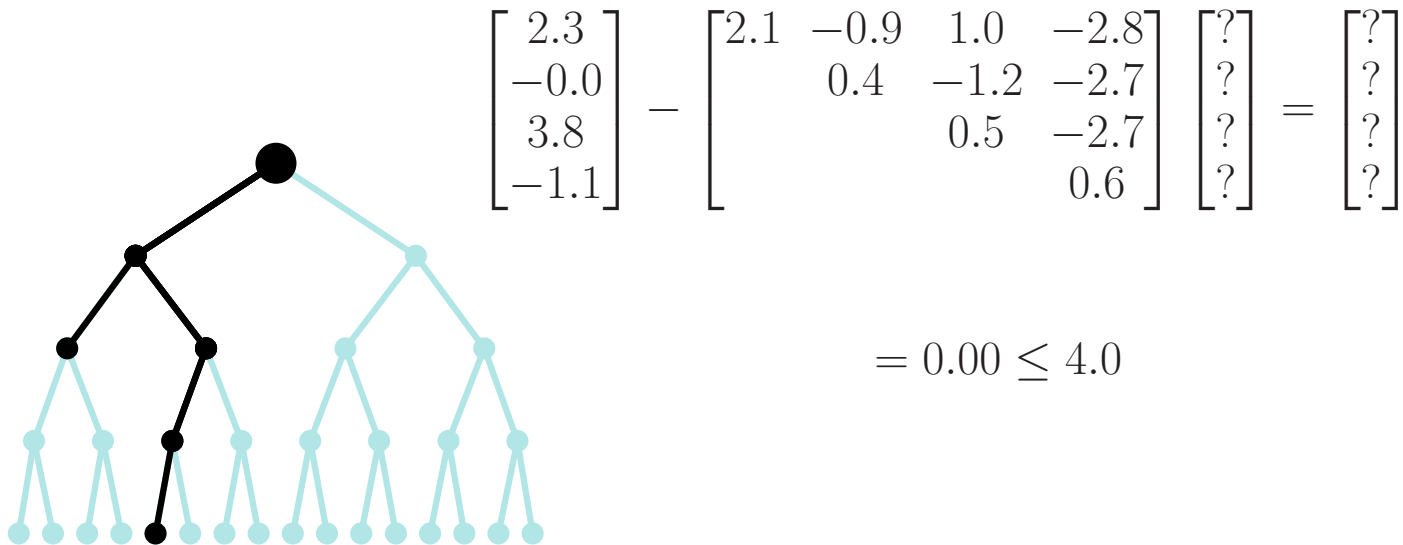
Sphere decoding example, numerics



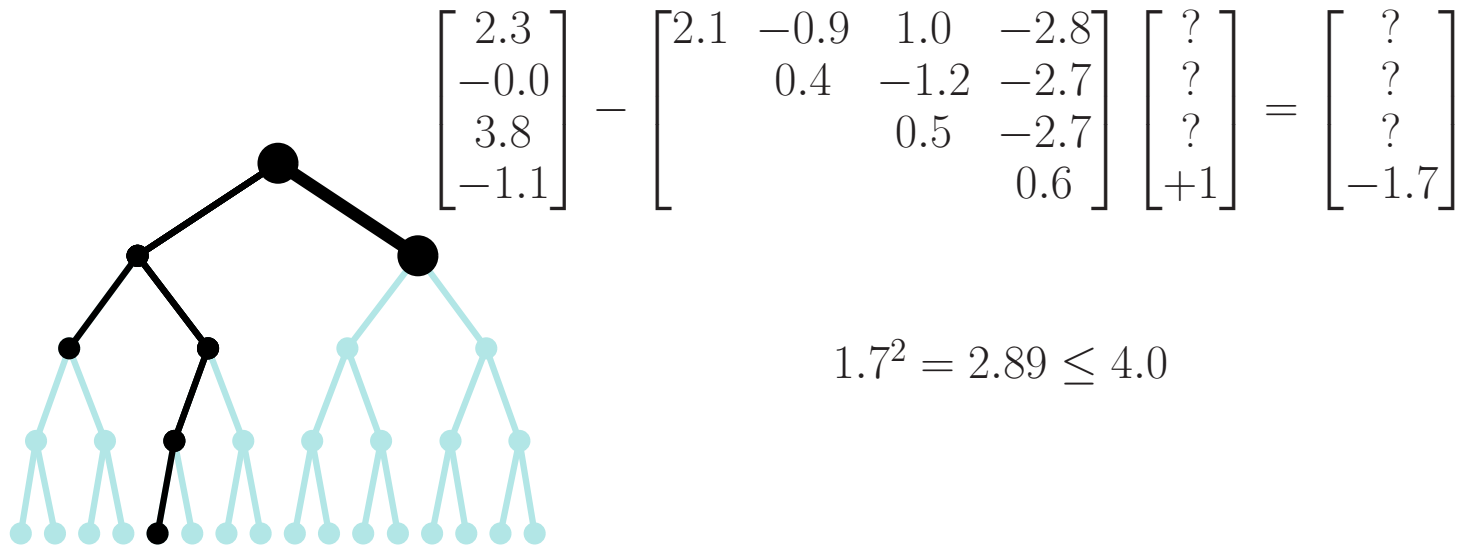
Sphere decoding example, numerics



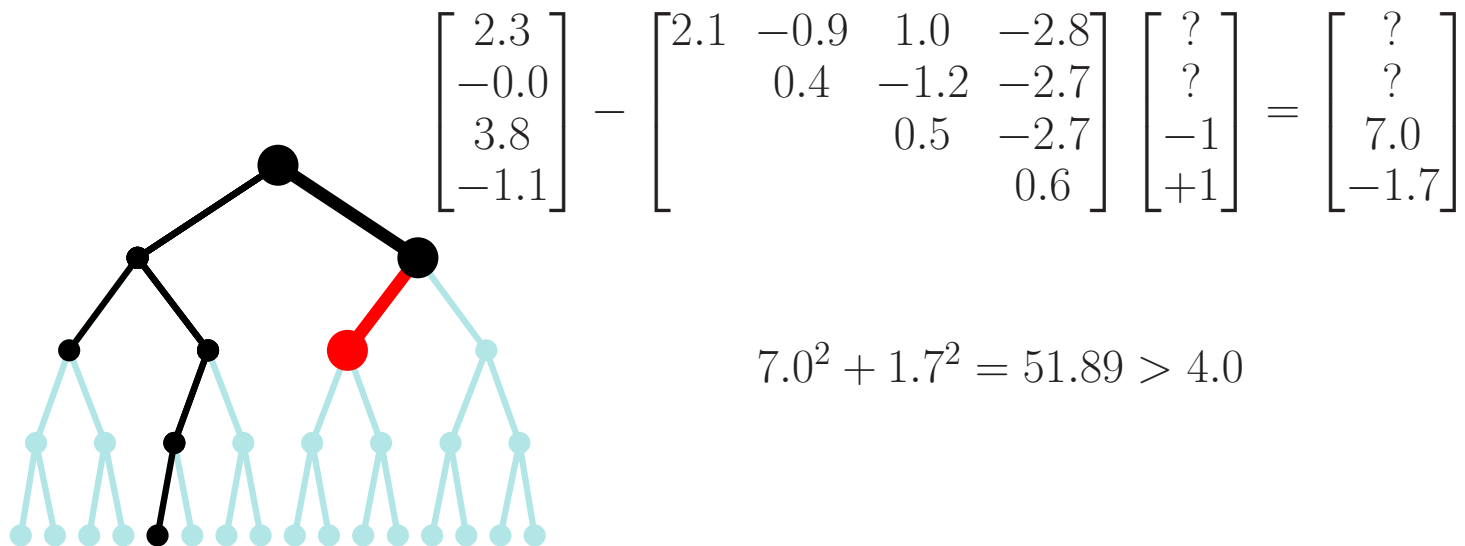
Sphere decoding example, numerics



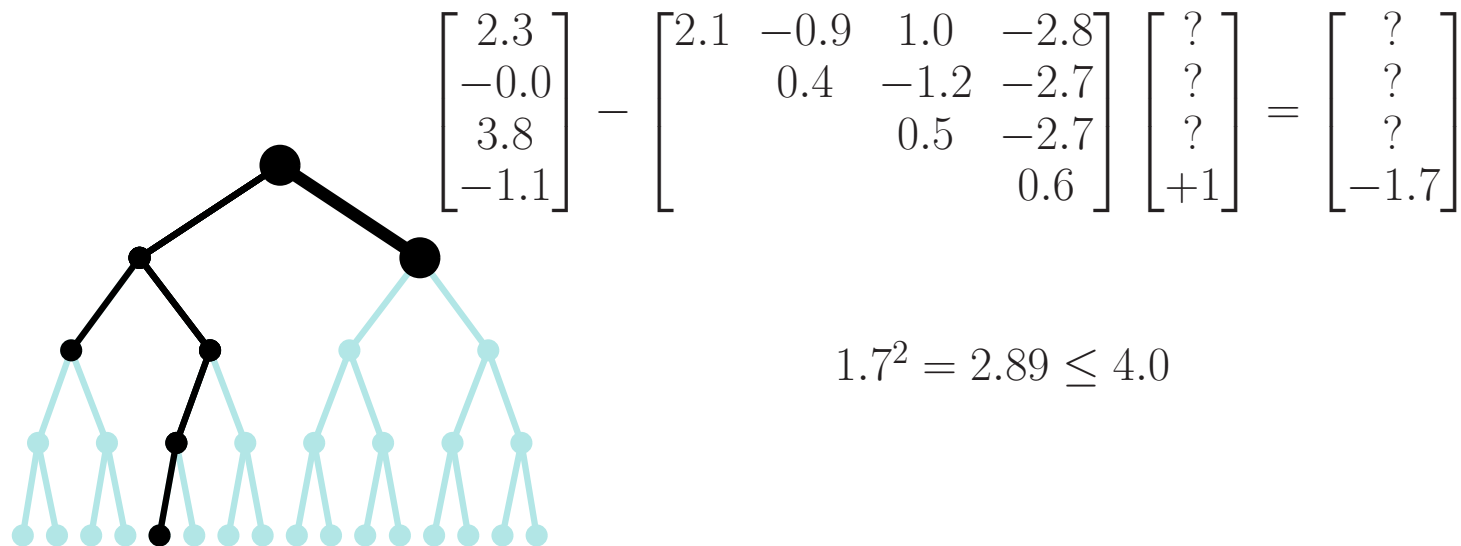
Sphere decoding example, numerics



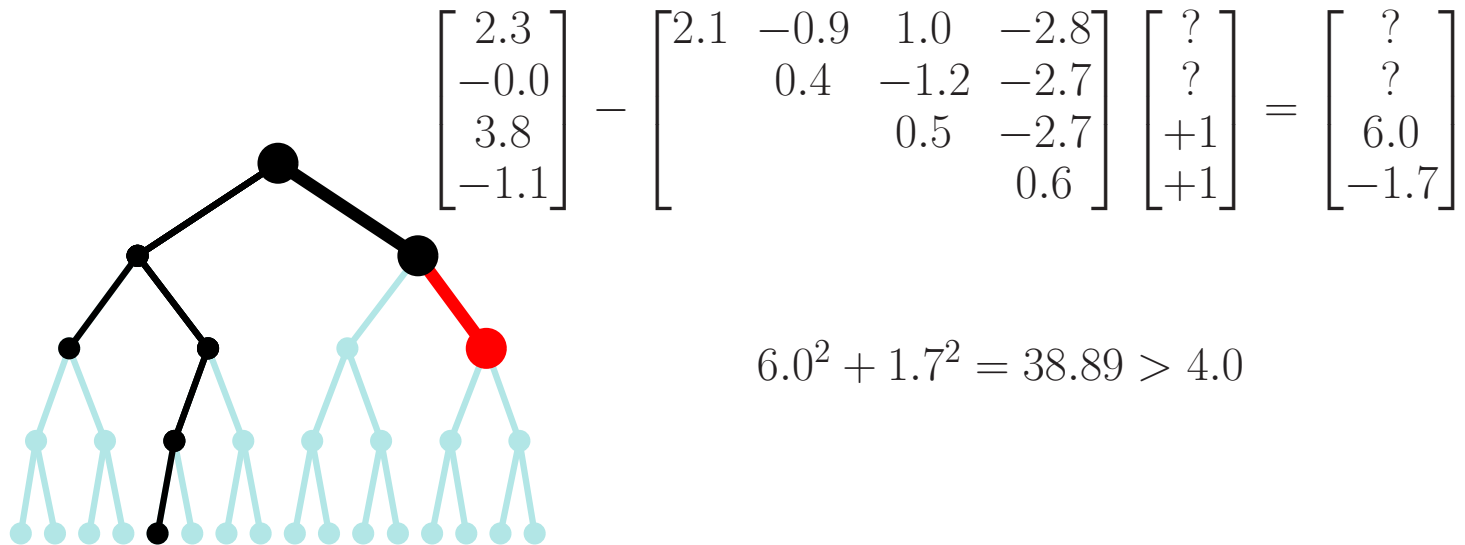
Sphere decoding example, numerics



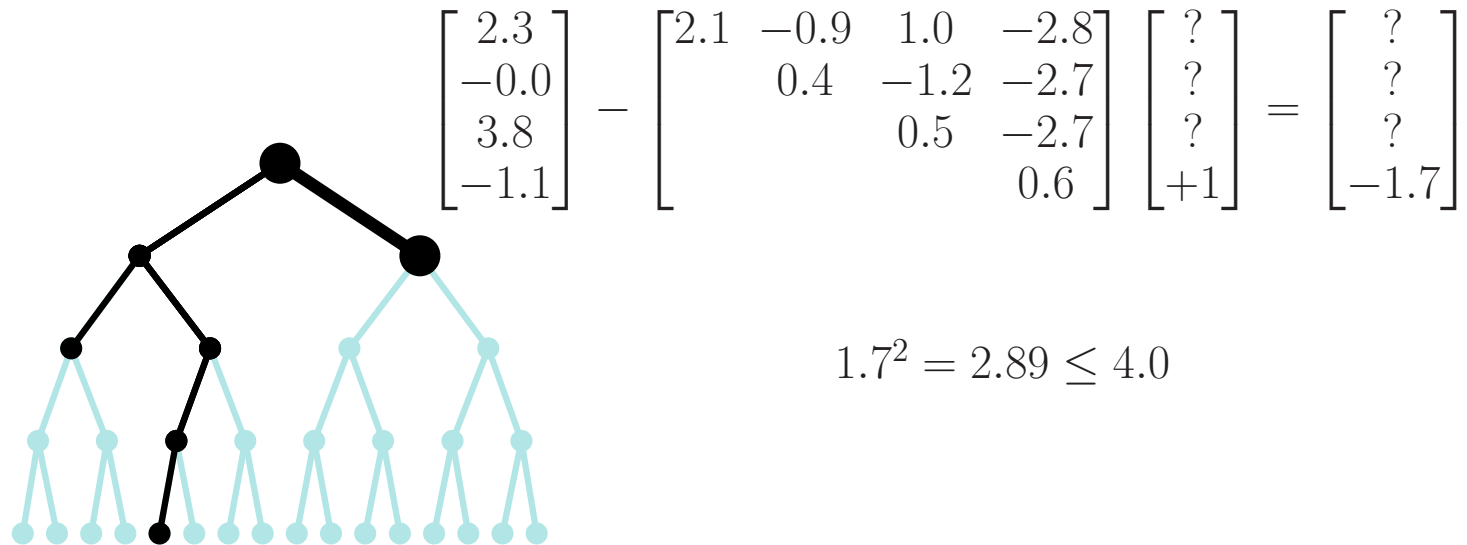
Sphere decoding example, numerics



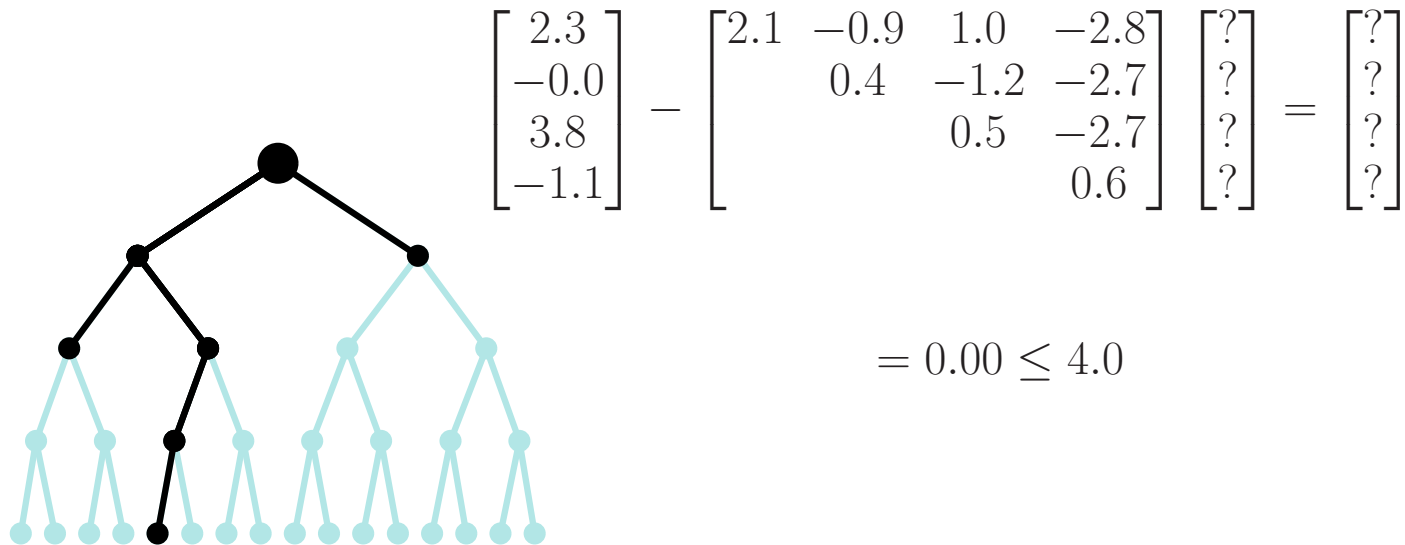
Sphere decoding example, numerics



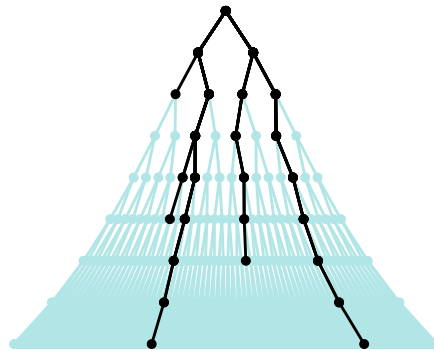
Sphere decoding example, numerics



Sphere decoding example, numerics



Accumulated complexity



Snapshot of instantaneous complexity

Universal bounds and equivalence of ML and lattice decoding

- We derived universal bounds for general MIMO (all scenarios, statistics, etc)
- We derived tightness whenever possible (broad setting)
- We will not get into that now: we focus on simpler more insightful settings

Theorem: *(Equivalence of ML and lattice decoding - Restatement)*
ML based sphere decoding and regularized lattice sphere decoding share the same complexity exponent for a very broad setting (share bounds and ‘tightness’)

⇒ ALL FOLLOWING RESULTS WILL HOLD FOR ML AS WELL AS FOR
(REGULARIZED) LATTICE SPHERE DECODING

UNIVERSAL BOUNDS - QUASI STATIC

Theorem: $c(r)$ is upper bounded as (piecewise linear)

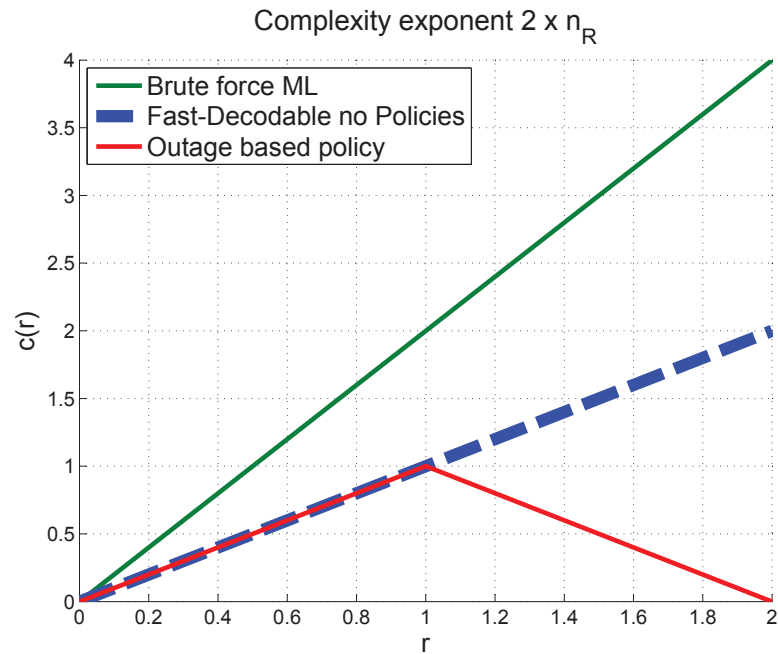
$$c(r) \leq \bar{c}(r) = \frac{T}{n_T} r(n_T - r), \quad r = 0, 1, \dots, n_T$$

for all fading statistics, all full rate lattice designs, and all decoding order policies.

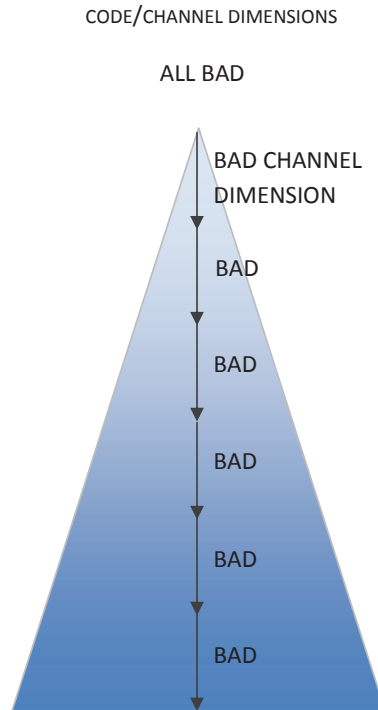
Example: $2 \times n_R$ Quasi-static

EXAMPLE: ($2 \times n_R$ channel ($n_R \geq 2$))

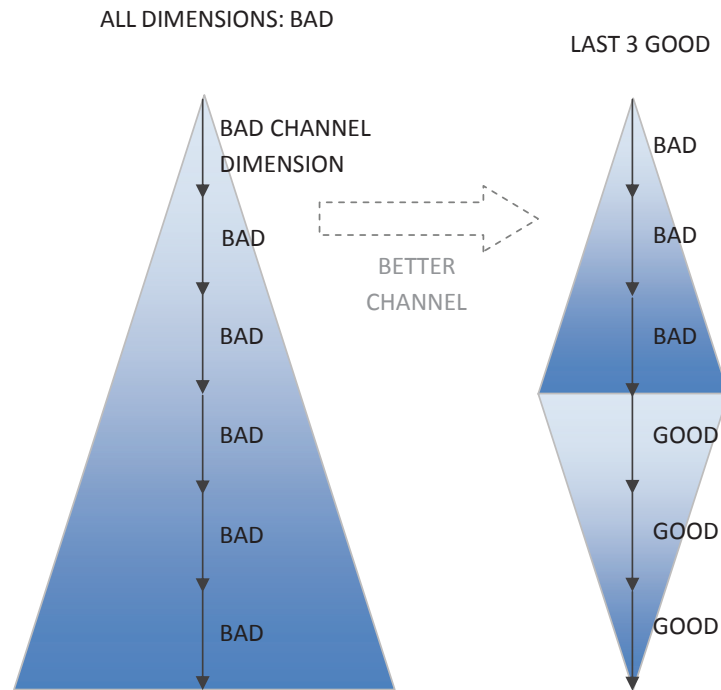
$$\bar{c}(r) = \begin{cases} r & r \leq 1, \\ 2 - r & r \geq 1. \end{cases}$$



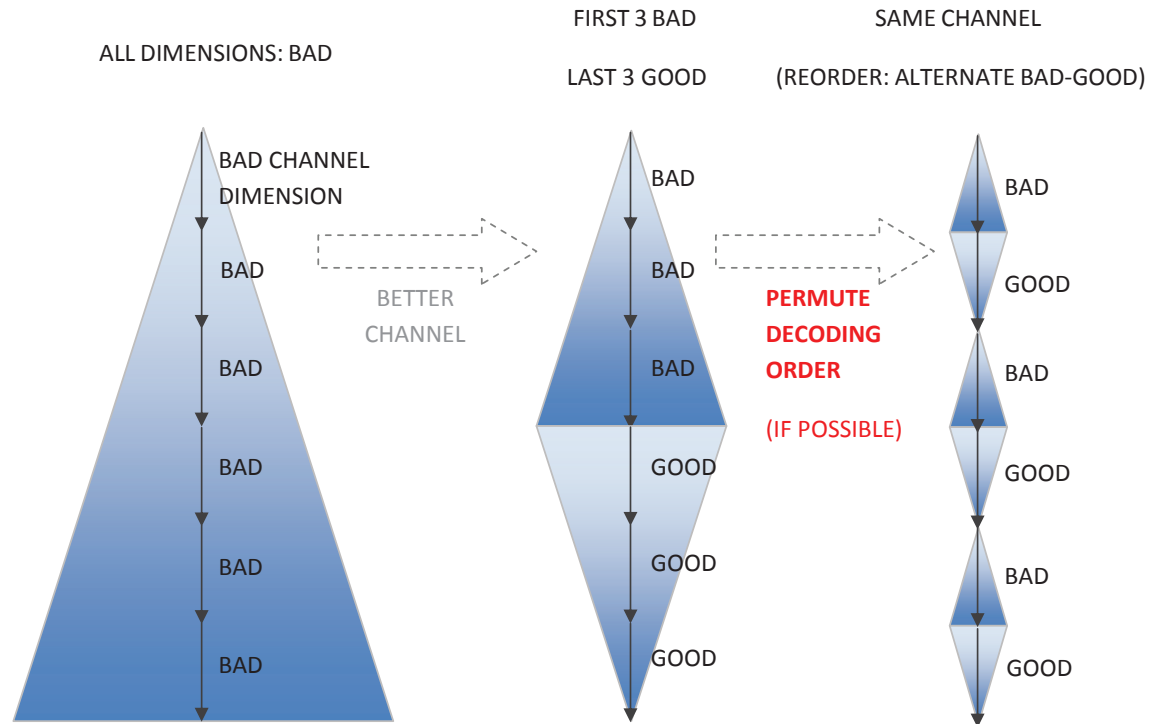
Decoding ordering and computational-halting policies



Decoding ordering and computational-halting policies₁



Decoding ordering and computational-halting policies₂



Tightness: DMT-opt quasi-static

TIGHTNESS OF UNIVERSAL BOUND

Theorem: (Quasi-static, Rayleigh, $n_R \geq n_T$) With probability 1 in the choice of the DMT optimal lattice design, the above is tight for all ordering policies.

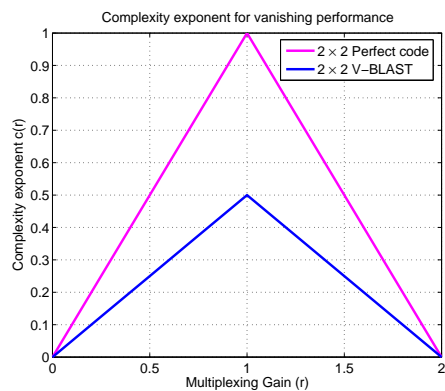
Theorem: (Quasi-static, Rayleigh, $n_R \geq n_T$) The bound is tight for all layered designs, for several fixed orderings including the natural ordering.

(Some hope remains for complexity reductions using dynamic policies)

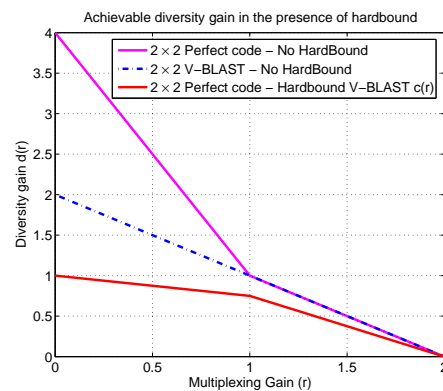
Complexity-constrained DMT

FIND MAX $d(r)$ GIVEN COMPLEXITY CONSTRAINT $N_{\max} \doteq \rho^{c_{\mathcal{D}}(r)}$ FLOPS

From uncoded to coded without increasing resources: a good idea?



(a) Complexity Hardbound



(b) Achievable DMT

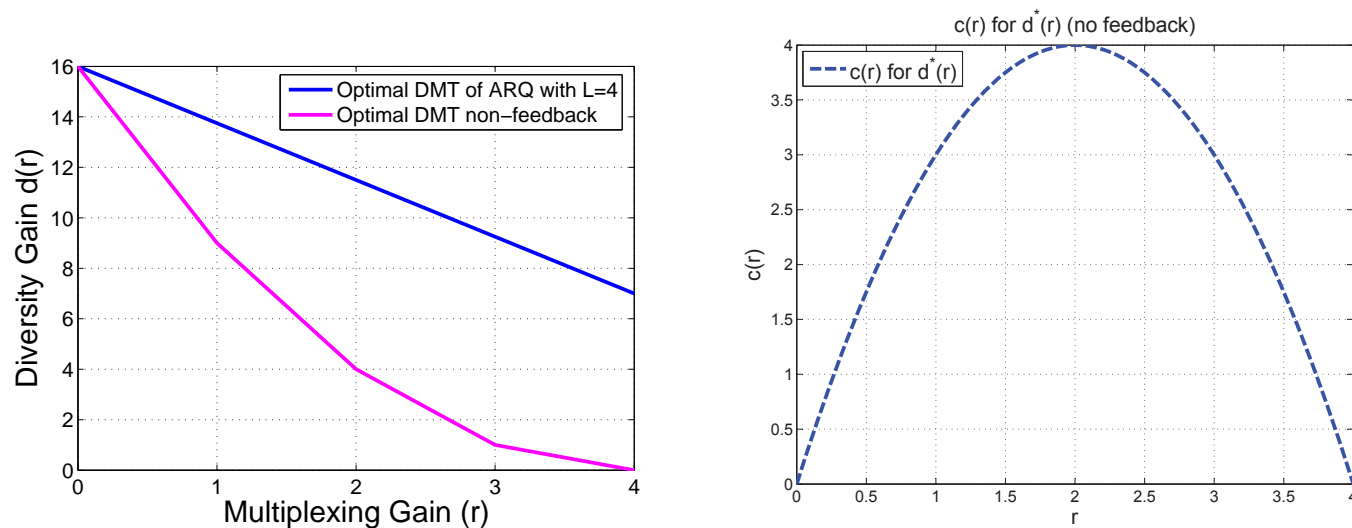
Performance-Complexity ramifications of feedback

PERFORMANCE-COMPLEXITY RAMIFICATIONS OF FEEDBACK

Performance-Complexity ramifications of feedback

Two interesting questions:

- What is the feedback-aided complexity to achieve DMT $d^*(r)$?
- What is the complexity to achieve the feedback-aided DMT¹ $d^*(r/L)$?



¹El Gamal, Caire, Damen

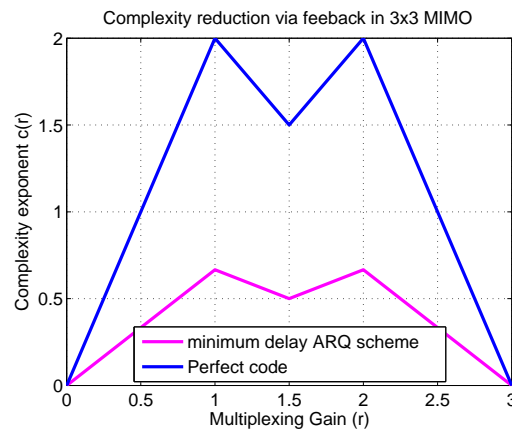
Feedback-aided complexity for optimal DMT $d^*(r)$

Corollary: (quasi-static iid Regular $n_R \geq n_T$, $LT = n_T$)

Minimum $c(r)$ for $d^*(r)$, (minimized over all lattice designs, all L -round ARQ schemes, all halting and decoding order policies), bounded as (piecewise linear $r = 0, 1, \dots, n_T$)

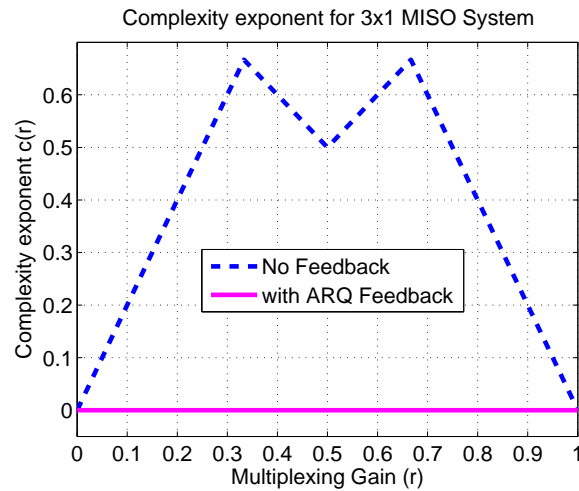
$$c(r) \leq \bar{c}_{red}(r) = \frac{1}{n_T} r(n_T - r).$$

- Compare to $c(r) = r(n_T - r)$
- Important role of “aggressive intermediate halting policies”



Feedback aided complexity for optimal DMT ($n_R < n_T$)

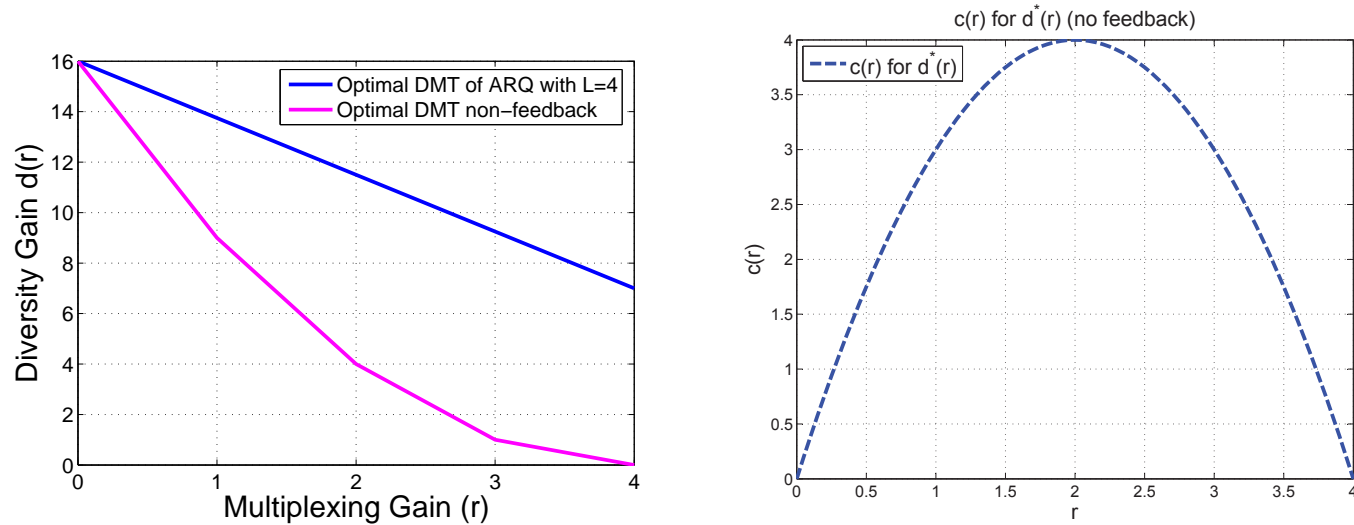
Corollary: $n_T \times 1$ MISO $L = n_T$, then $c(r) = 0$ for $d^*(r)$



Complexity cost for feedback-aided DMT $d^*(r/L)$

SEEKING TO $c(r)$ NEEDED TO ACHIEVE $d^*(r/L)$

Recall:



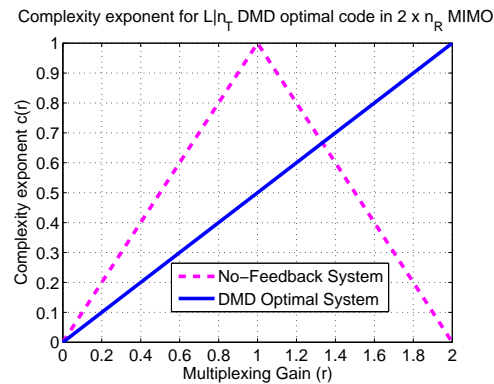
Complexity reduces with feedback despite increased $d^*(r/L)$

Theorem: ($L|n_T$, quasi-static, $n_R \geq n_T$)
 Minimum $c(r)$ to achieve optimal $d^*(r/L)$ is bounded as ((mult. of L))

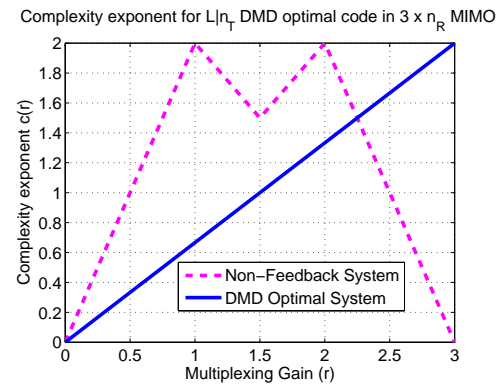
$$c(r) \leq \bar{c}_{dmd}(r) = \frac{rn_T}{L^2} \left(L - \frac{r}{n_T} \right).$$

Corollary: The above with $L = n_T$ gives

$$c(r) \leq \bar{c}_{DMD}(r) = \left(1 - \frac{1}{n_T} \right) r.$$



$2 \times n_R$

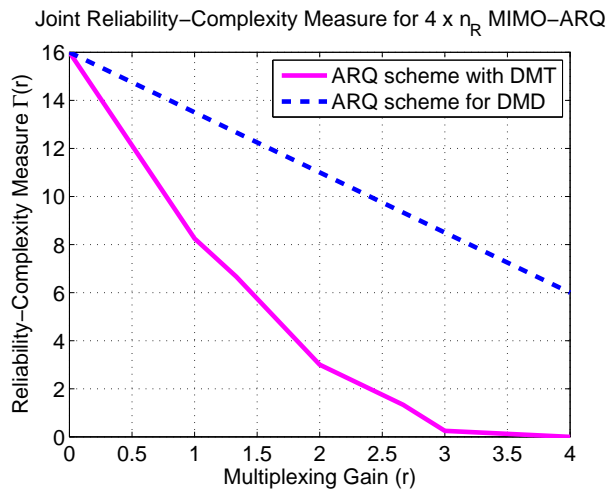


$3 \times n_R$

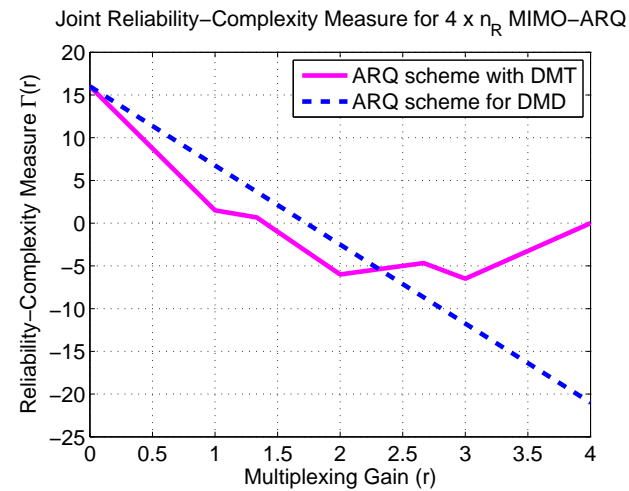
Have feedback: Go for basic DMT or feedback-aided DMT?

Joint performance-complexity measure

$$\Gamma(r) = d(r) - \gamma c(r)$$



$\gamma = 1$



$\gamma = 10$

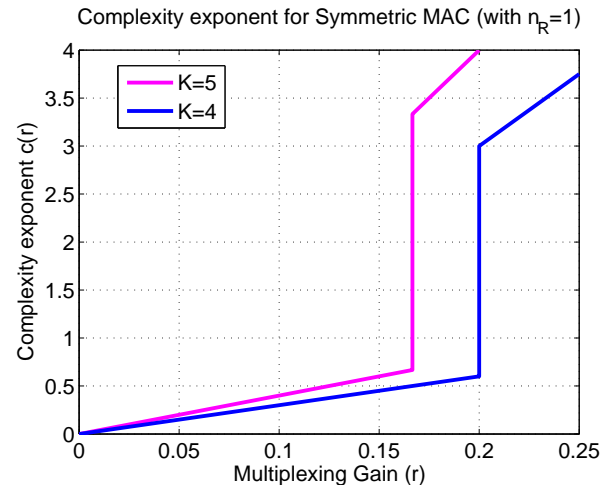
Joint reliability-complexity measure for DMT and DMD optimal ARQ schemes

Multiple access channel

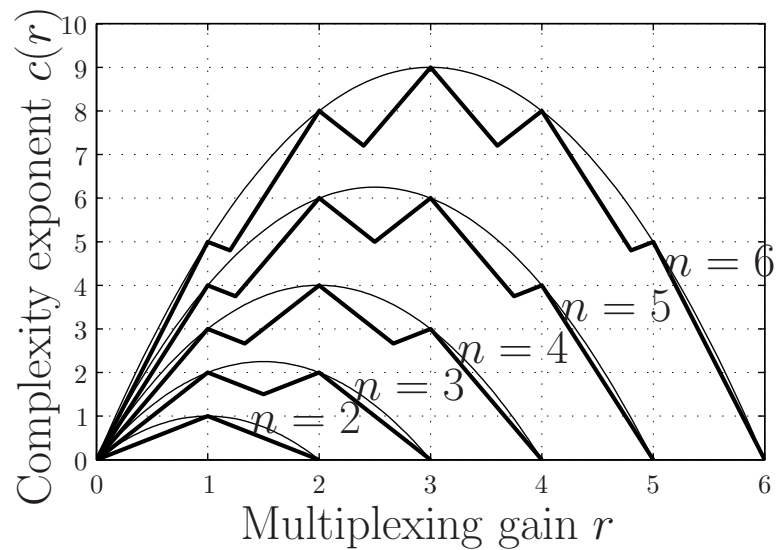
MIMO MAC

Corollary: (K user MAC, $n_T = 1$, $n_R = 1$, r per user, Rayleigh, K odd)
The minimum $c(r)$ (over all lattice designs and halting and decoding order policies) to achieve the optimal MAC-DMT, is upper bounded as

$$c(r) \leq \bar{c}_{mac}(r) = \begin{cases} (K-1)r & \text{for } r \leq \frac{1}{K+1}, \\ (K-1)Kr & \text{for } \frac{1}{K+1} < r \leq \frac{1}{K}. \end{cases}$$



Conclusion on search based algorithms



CONCLUSION FOR ML AND (REGULARIZED) LATTICE BASED SOLUTIONS

- Very considerable complexity for high performance
 - ★ Feedback helps
- No known way to drop below the upper bounds

Lattice reduction

LETS GET SOME HELP FROM LATTICE REDUCTION (LR)

BUT REMEMBER

LR PROBLEMATIC IN UNAVOIDABLE SCENARIOS (INNER-OUTER CODE)

Lattice reduction

CHANGE REPRESENTATION OF LATTICE

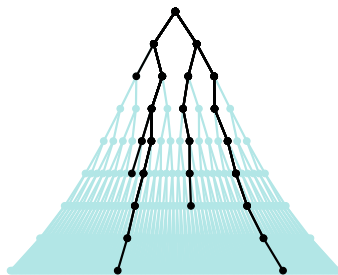
LR: Input \mathbf{H} , Output \mathbf{T} (unimodular matrix of integers)

$$\underbrace{\mathbf{H}}_{\text{CHANNEL}} \cdot \underbrace{\mathbf{s}}_{\in \mathbb{Z}} = \underbrace{\mathbf{HT}^{-1}}_{\text{better channel}} \cdot \underbrace{\mathbf{T}\mathbf{s}}_{\text{still } \in \mathbb{Z}},$$

Achieving optimal exponents - but unbounded gap

Theorem: (Trans-IT Oct 2010) *LR-aided regularized linear decoding and LR-based halting, achieves $c(r) = 0, d(r) = d_{opt}(r)$, for all r , all codes, all MIMO scenarios and all fading statistics. (at most $O(n^2)$ flops per bit)*

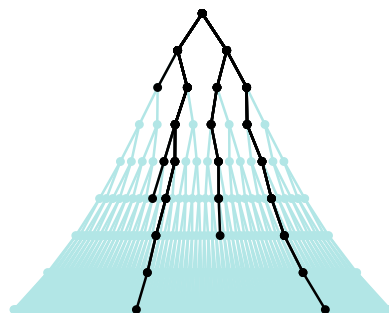
- First ever solution to achieve optimal $d^*(r)$ with subexponential complexity.



- BUT! Potentially unbounded gap to exact lattice decoding!

Achieving a vanishing gap at subexponential complexity

Theorem: (Trans-IT subm. July 2011) LR-aided regularized lattice sphere decoding with LR- and outage-based halting policies, introduces a zero complexity exponent, and achieves a vanishing gap to the exact implementation of lattice decoding (all MIMO scenarios, all statistics, all codes).



- First ever to achieve a vanishing gap to the exact solution of (regularized) lattice decoding, with subexponential computational complexity
- Again though - remember LR limitation!

Conclusions: Flops for Ergodicity

Can small chips (rather than CSIT) give us ergodicity?

- With LR - Yes: for a very broad setting, but not for near-ergodic rates
- With LR and 1 bit of feedback - Yes: for all r
- BUT: LR might not apply
- Without LR: Mostly NO - open problem - there might be hope!
 - ★ A little bit of feedback goes a long way
 - ★ High multiplexing gain most problematic

Other contributions

OTHER RESEARCH CONTRIBUTIONS

Other contributions

- Cooperative wireless networks (Trans. IT 2009)+(subm. Trans-IT 2011)
- Two-way multi-directional communications
 - ★ patents (pending) - publications - funding - award
- Cross layer optimization - queue/channel (Trans. IT 2009)
 - ★ Towards a consummated union: finite delay results
- Feedback (Trans. IT 2009)
- Connectivity in networks with bounding constraints (publications)

Other contributions₁

- Soft-biometrics / surveillance networks / computer vision (publications)²
 - ★ Will spend some time at well known biometrics lab in US
- Interference (preliminary work)
 - ★ Interference alignment and diversity (potential submission ISIT-2012)
 - ★ Stale CSIT (potential submission ISIT-2012)
 - ★ Uplink-downlink DOF (ITA-2011, ISIT-2011)
 - ★ Finite SNR IA. LR-aided IA.

²“Search pruning video surveillance systems: Efficiency-reliability tradeoff,” 1st IEEE Workshop on Information Theory in Computer Vision and Pattern Recognition, Nov. 2011.

Other than that

- Successful funding efforts (e.g. ANR Blanc International)(also on surveillance)
- Relatively close to industry

- Tutorials + awards (for last two results + two-way)
- Efforts to recruit talented students

Thank you

THANK YOU VERY MUCH FOR COMING