

Security in Wireless Ad hoc Networks

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Mobile Ad Hoc Networks (MANET)

- Collection of wireless mobile hosts forming a temporary network
- No fixed network infrastructure
- No (or limited) organization

- Military and Emergency
- Sensor Networks
- Civilian applications, ubiquitous computing

Trust in MANET

- Managed environment
 - A-priori trust
 - Entity authentication \Rightarrow correct operation
 - But:
requirement for authentication infrastructure
- Open environment
 - No a-priori trust
 - authentication does not guarantee correct operation
 - *New security paradigm*

Node Misbehavior

Selfish Nodes

- Do not cooperate
- Priority: battery saving
- No intentional damage to other nodes.
- **Exposure:**
 - **passive** denial of service
 - black hole
 - idle status

Malicious Nodes

- Goal: damage to other nodes
- Battery saving is not a priority
- **Exposure:**
 - **active attacks**
 - denial of service
 - traffic subversion
 - attacks exploiting the security mechanism

MANET Requirements

Wireless & Mobile

- Limited Energy
- Lack of physical security

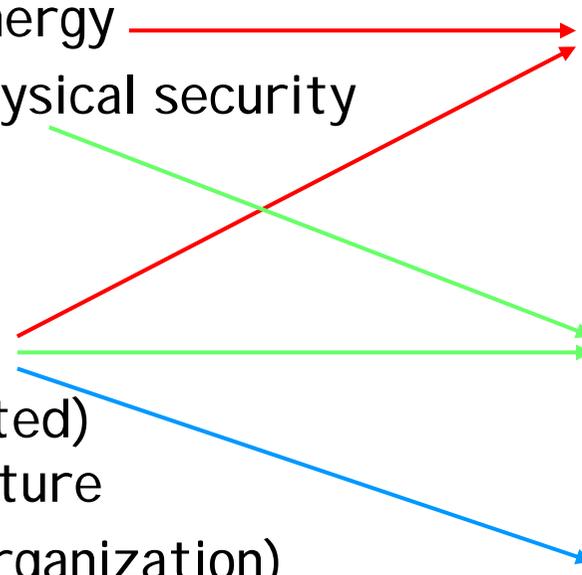
Ad Hoc

- No(or limited) infrastructure
- (Lack of organization)

• Cooperation Enforcement

• Secure Routing

• Key management



Cooperation Enforcement in MANET

- Routing and Packet Forwarding cost energy.
- Selfish node saves energy for itself
- Without any incentive for cooperation network performance can be severely degraded.

[Michiardi, Molva EW'02]

Cooperation enforcement mechanisms

Token-based [Yang, Meng, Lu]



Threshold cryptography

Nuglets [Buttayan, Hubaux]

SPRITE [Zhong, Chen, Yang]



Micro-payment

CONFIDANT [Buehgger, Le Boudec]

CORE [Michiardi, Molva]

Beta-Reputation [Josang, Ismail]

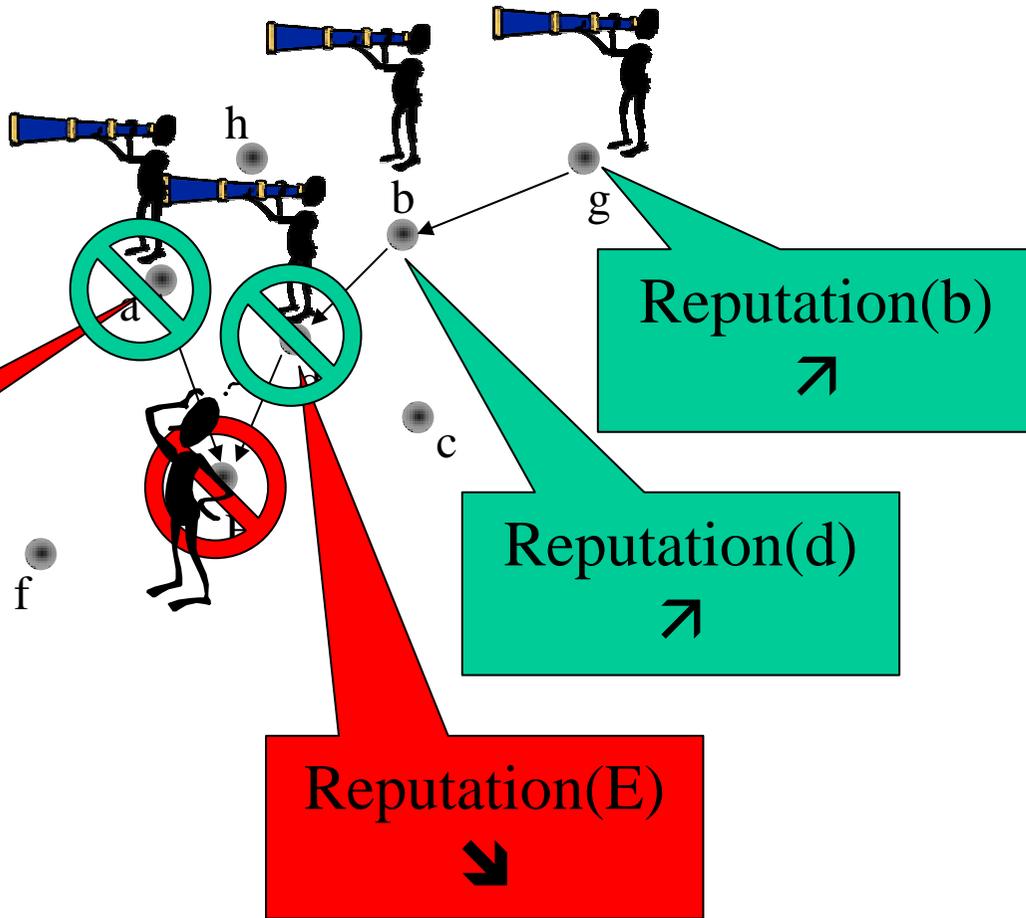


Reputation-based

CORE

Packet forwarding
Source Node: **g**
Destination Node: **f**
Route: $\langle g, E, d, E, f \rangle$

Reputation(E)
↓



Reputation(E)
↓

Cooperation Enforcement Evaluation with Game Theory

- Cooperative GT
 - Study the *size* (k) of a *coalition* of cooperating nodes

utility function : $U(k) = \alpha_i u(y_i) + \beta_i r(\sigma_i)$

relative share : $\sigma_i = \frac{y_i}{\sum_j y_j}$

- Nash Equilibrium \rightarrow lower bound on k

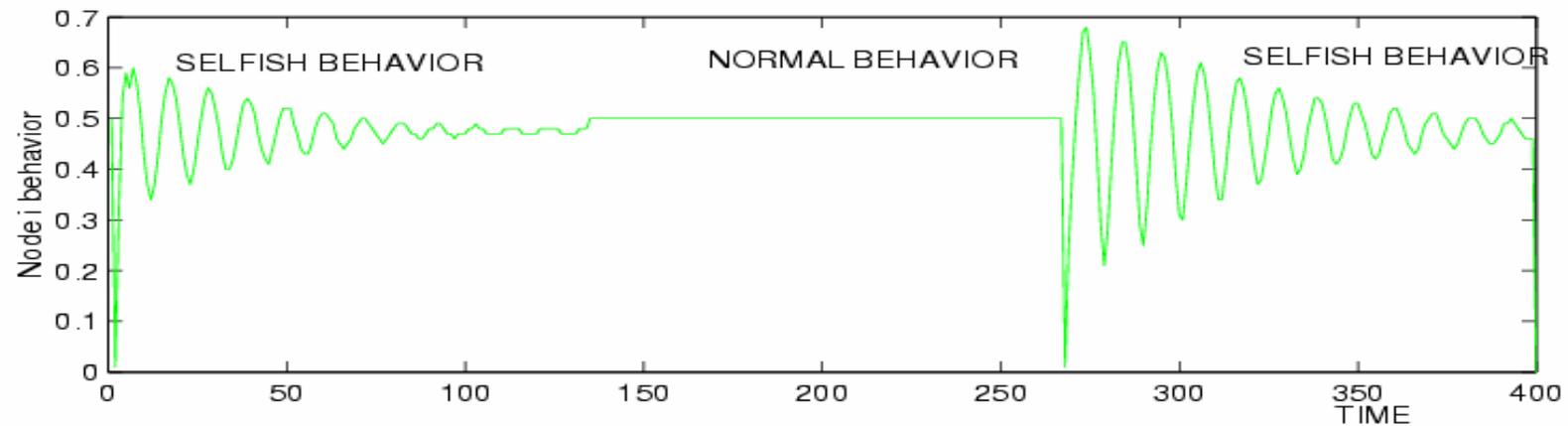
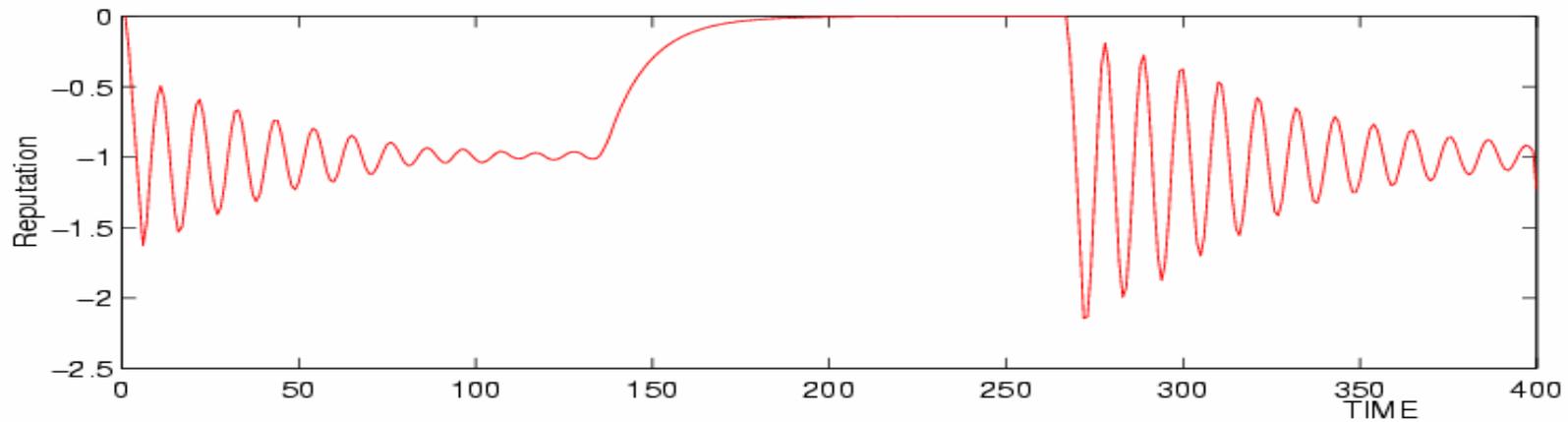
- Non-cooperative GT

- Utility function with *pricing*

$$u_i(b_i, b_j) = f(E_{self}, E_R, E_{PF}, b_i, b_j, r_i)$$

- Pricing used to guide the operating point (i.e. maximum of utility function) to a fair position
- r_i : dynamic reputation of node n_i evaluated by her neighbors

Non-cooperative GT with pricing



Secure Routing - Vulnerabilities

- Modification
- Impersonation
- Fabrication
- Wormhole attack
- Lack of cooperation

Secure Routing - Objectives

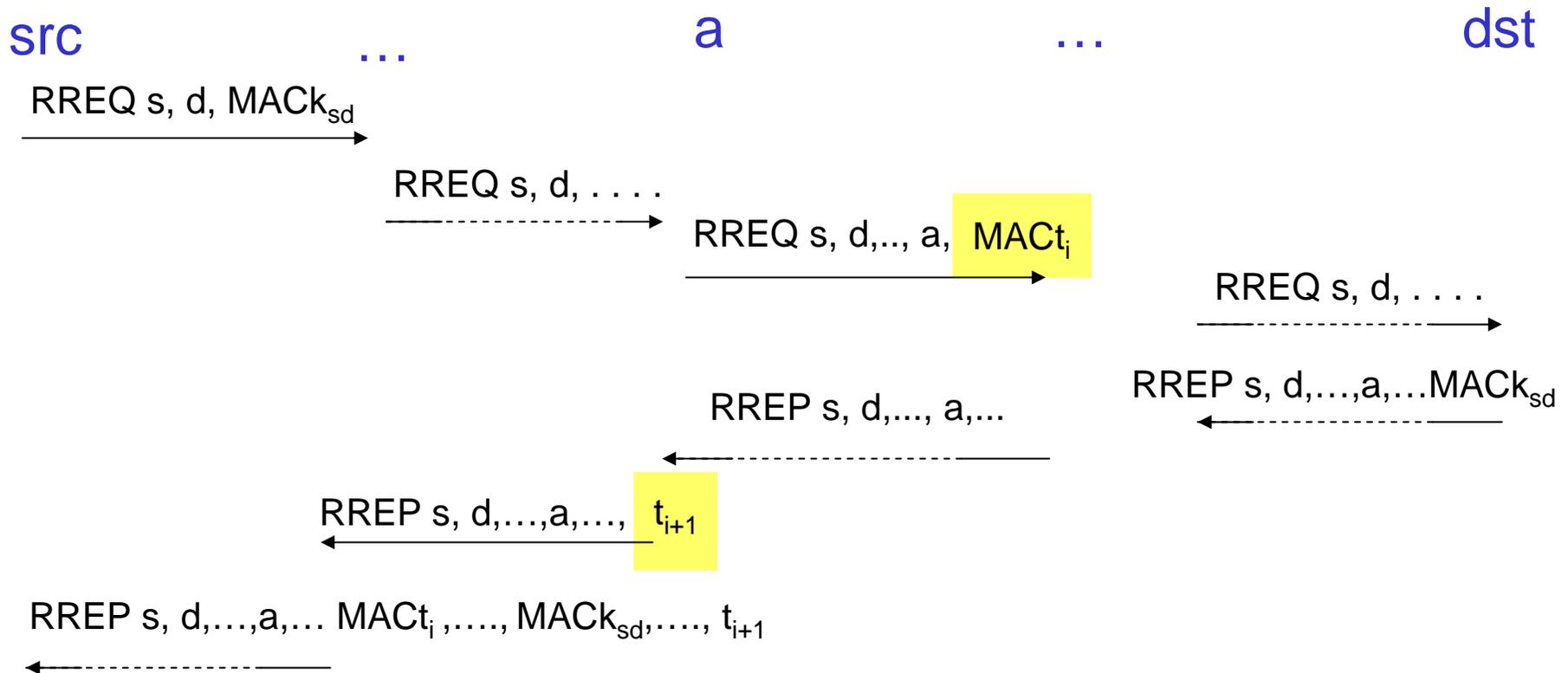
- Authentication (Integrity) of routing information
- Entity authentication
 - Source
 - Destination
 - Intermediate node
- Correct behavior (of algorithm, if any)
- Asymmetric vs. Symmetric Crypto
- Pro-active vs. Reactive routing protocols

Routing in MANET

- Reactive (on-demand)
 - Dynamic Source Routing (DSR)
 - Ad Hoc On-demand Distance Vector (AODV)
- Pro-active
 - Destination Sequenced Distance Vector (DSDV)
 - Optimized Link State Routing (OLSR)
- Hybrid
 - Zone Routing Protocol (ZRP)
 - Distributed Dynamic Routing (DDR)
- Location-based
 - Location-Aided Routing (LAR)

ARIADNE [Hu, et al.]

- On-demand Routing Protocol DSR
- k_{sd} : shared secret known by (src, dst)
- $t_i = h^{n-i}(\text{secret})$: TESLA key of a node valid for time interval T_i
disclosing t_{i+1} in T_{i+1} authenticates the node



Prerequisite: distribution of authenticated TESLA keys ($h^n(\text{secret})$)

Other Secure Routing Proposals for MANET

- Secure Routing Protocol [Papadimitriou, Haas]
 - security associations between source end destination only
- ARAN [Dahill, et al.]
 - PK certificates for IP @
- SEAD [HU, et al.]
 - proactive routing authenticated hash chains
- TESLA with instant key disclosure (TIK)
 - can cope with wormhole attack

Secure Routing Summary

- No new requirement other than self-organized key management
- All solutions rely on some key set-up prior to secure routing operation
- Contradiction: long-lived security associations in self-organized MANET

Key Management Requirements

- Secure routing
- Basic security services
 - Authentication
 - Confidentiality
 - Integrity
 - Non-repudiation
- Symmetric or Asymmetric Keys

Key Management Challenges

Lack of (or limited)

- Security infrastructure
 - Key servers (KDC, CA, RA)
- Organization (a priori trust)
 - p2p
 - Authentication is not sufficient to build trust

Key Management Objectives

- Bootstrapping from scratch
- Fully distributed
- Minimum dependency

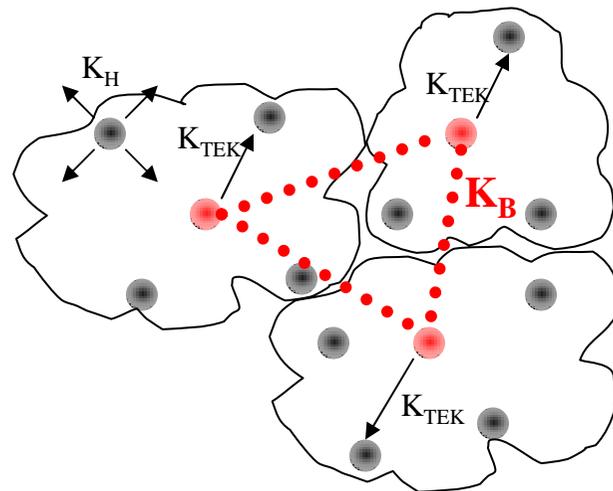
Key Management Approaches

- Based on symmetric crypto
- (I D, PK) binding
 - PK Certificate = $(I D, PK)_{CA}$
 - Self-organized CA
 - Web of trust(PGP)
 - No certificate
 - Crypto-based I Ds: $I D = h(PK)$
 - I D-based Crypto: $PK = f(I D)$
- Context-dependent authentication
 - location-limited channels
 - Shared passwords

Key Management Based on Symmetric Cryptography

Secure Pebblenets [Basagni et al.]

- cluster formation algorithm



- Cluster Head
- Cluster Member

K_G = group key, **well known**

K_H = hello key (derived from K_G), used for cluster head selection

K_B = inter cluster key, used for traffic encryption key generation

K_{TEK} = used for traffic confidentiality

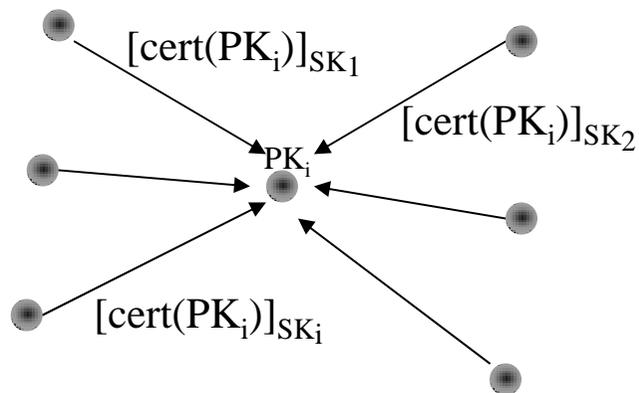
Assumption: no malicious nodes

(ID, PK) binding

Self-organized CA

[Zhou, Haas] [Kong, et al.] [Yi, Kravets] [Lehane, et al.]

- Based on threshold cryptography



$CERT(PK_i)_{SK}$ {
[cert(PK_i)]_{SK₁}
[cert(PK_i)]_{SK₂}
...
[cert(PK_i)]_{SK_i}
...}

Verification of $CERT(PK_i)_{SK}$ by any node
using well known PK

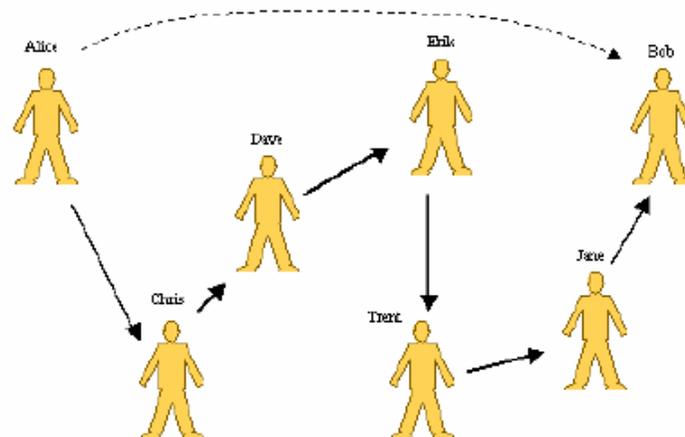
- PROs: distributed, self-organized
- CONs: share distribution during bootstrap phase, network density, Sybil attack

(ID, PK) binding

Web of Trust (PGP)

[Hubaux, Buttyan, Capkun]

- No CA
- Alice → Bob and Bob → Eve ⇒ Alice → Eve
- Merging of certificate repositories



- PROs: no centralized TTP
- CONs: initialization, storage, transitivity of trust

(ID, PK) binding

Crypto-based ID

- SPKI [Rivest]
- Statistically Unique Cryptographically Verifiable IDs [O'Shea, Roe] [Montenegro, Castellucia]
IPv6 @ = NW Prefix | h(PK)
- DSR using SUCV-based IP addresses
[Bobba, et al]

PROs: no certificates, no CA

CONs: generation of bogus IDs

(I D, PK) binding

I D-based Crypto

[Halili, Katz, Arbaugh]

[Boneh, Franklin, CRYPTO 2001]

- I D-based
 - $PK = h(I D)$
 - SK computed by TTP
- Threshold Crypto to distribute TTP

PROs: no certificates, no centralized server

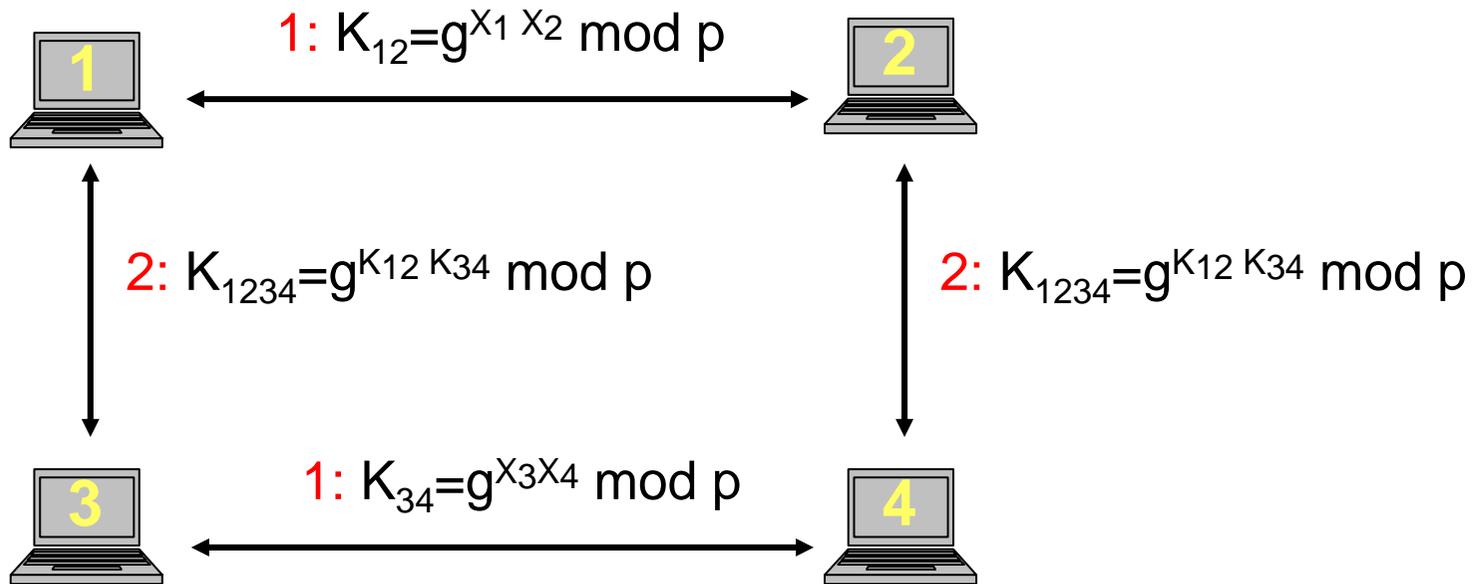
CONs: distribution of initial shares

Context-dependent Authentication

Password Authenticated Key Exchange

[Asokan, Ginzborg]

HyperCube Protocol (Diffie-Hellman)



PROs: self-organized, fully distributed

CONS: shared password

Context-dependent Authentication

Secure channel

- [Balfanz, et al.] establish pairwise security association based on vicinity of devices
- [Capkun, et al.] secure channel + web of trust
- PROs: self-organized, fully distributed
- CONs: reliance on secure side channel

Layer 2 vs MANET Security

- IEEE 802.11 and Bluetooth
 - weaknesses
 - secure extensions to wireline networks
- Layer 2 mechanisms in MANET
 - managed environments: L2 sufficient if node integrity is guaranteed (tamper-proof HW)
 - open environments (no a priori trust): L2 cannot cover higher layer (3,4, ..) security

State of the art - Summary

- Specific requirements
 - Cooperation enforcement
 - Bootstrapping security associations
- Solutions yet to come . . .
- Interesting applications of cryptography
- Some untruths and non-sense

Main Flaw

- Security requirements in MANET are stronger than in “classical” networks.
- MANET networking still is a research topic
- Security retrofitted as add-on mechanisms as if network technology was established.

Right Approach

- Address security at early stages of protocol design: i.e. Routing Protocol dealing with Routing+Cooperation+Key Management
- Old model based on verification of credentials and authentication not suitable, identities are meaningless
- Further develop & integrate new concepts
 - *A posteriori* trust (based on observation, reputation, imprinting)
 - Partial assurance
 - Substitute infrastructure with context information (location, physical distance, history)
 - Others to be invented

Conclusion

- Wireless Ad Hoc Security still in its infancy
 - Lack of integrated approach
 - Looking for suitable new paradigms
 - Partial coverage (privacy, intrusion detection, physical attacks, etc.)
- ⇒ Room for creativity

THANK YOU