

Near-far gain in a multiuser diversity system

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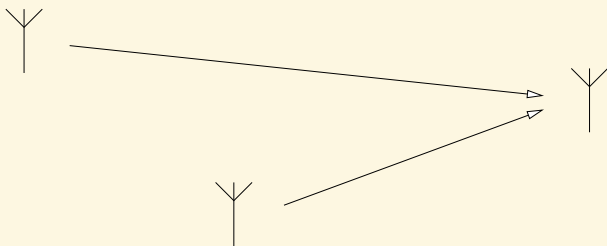
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- Time varying channels
- Unbalance between channels
- Selection diversity over users



Sum-rate greedy scheduler

Schedule one user with strongest channel \rightarrow time division.

Hard-fair scheduler

Schedule all users (invert channel, guarantee rate) \rightarrow code division.

Initial thought

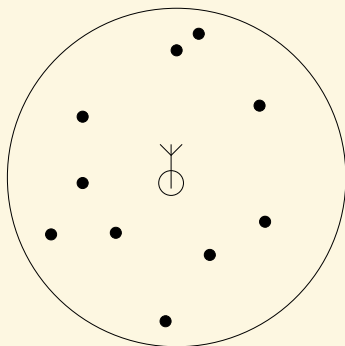
- PFS is sum-rate suboptimal
(maximizes $\sum_{k=1}^K \log(E[R_k])$).
- We should be able to improve by simultaneous scheduling

What shall we analyze

- Average spectral efficiency / sum rate.
- Focus on high and low spectral efficiency regimes
(wideband regime)

Single cell system – uplink (multiple access channel)

- K users.
- Each user has fixed random pathloss s_k (exponential).
- Each user has i.i.d. short term fading f_k (Rayleigh).
- Channel gain of user k is $d_k = s_k f_k$.



At time t PFS schedules the user k with maximum metric

$$\frac{R_k(t)}{T_k(t)}$$

where

$$R_k(t) = \log_2 \left(1 + \frac{d_k(t)}{N_0} \right) \quad (1)$$

$$T_k(t+1) = \lambda T_k(t) + (1 - \lambda) R_k(t) \quad (2)$$

The forgetting factor $\lambda \in [0, 1)$

if $\lambda \rightarrow 0$, PFS is more round-robin

if $\lambda \rightarrow 1$, PFS is at its greediest

When $\lambda \rightarrow 1$ (infinite forgetting), PFS

- schedules the user with largest short term fading f ,
- allocates rate according to channel state d ,
- enforces orthogonal transmission (TDMA),
- provides no rate guarantee, and
- is (still) not sum rate optimal.

Key idea

- Choose K_A users with largest short term fading f_k
- Allocate each user the power $1/K_A$
- Allocate each user rate assuming superposition coding.

Answers we have

- System capacity behavior.
- Optimal number of simultaneous users.

Answers to find

- Can we formulate a simple recursive allocation cf. PFS?
- What did we just do to “fairness”?

Scheduling $k \in \mathcal{A}$

- Received energy $y = \frac{1}{K_A} \sum_{k \in \mathcal{A}} f_k s_k$.
- Sum rate $R(y) = \log_2(1 + y\text{SNR})$.

Average system capacity and system energy per bit

$$C = \int_0^{\infty} R(y) dF(y) \quad (3)$$

$$\left(\frac{E_b}{N_0}\right)_{\text{sys}} = \frac{\text{SNR}}{C} \quad (4)$$

$F(y)$ is the distribution of the K_A largest f_k (out of K) each multiplied with a random path loss s_k .

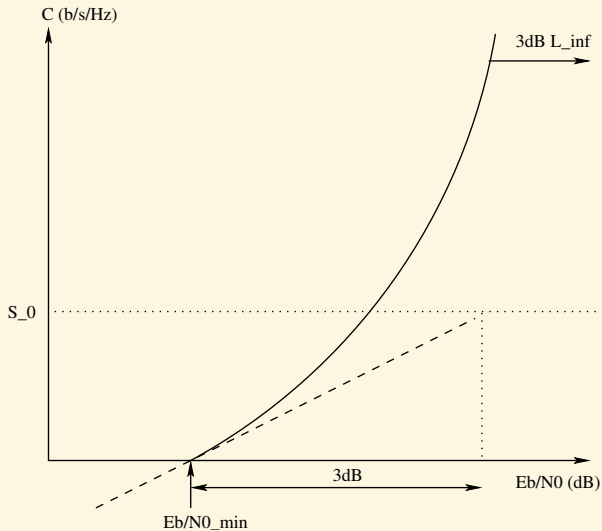
Low spectral efficiency approximation

$$\left(\frac{E_b}{N_0}\right)_{\text{sys}} \Big|_{\text{dB}} = \left(\frac{E_b}{N_0}\right)_{\text{min}} \Big|_{\text{dB}} + \frac{3\text{dB}}{\mathcal{L}_0} C + o(C), \quad (5)$$

High spectral efficiency approximation

$$\left(\frac{E_b}{N_0}\right)_{\text{sys}} \Big|_{\text{dB}} = \frac{3\text{dB}}{\mathcal{L}_\infty} C - 10\log_{10}(C) + \mathcal{L}_\infty 10\log_{10}(2) + o(1), \quad (6)$$

The same visually



$$\left(\frac{E_b}{N_0}\right)_{\min} = \frac{\log(2)}{C'(0)} \quad (7)$$

$$= \frac{\log(2)}{\int_0^{\infty} \frac{y}{1+\text{SNR}y} dF_y(y) \Big|_{\text{SNR}=0}} \quad (8)$$

$$= \frac{\log(2)}{E[y]} \quad (9)$$

$$= \frac{\log(2)}{E[s] \frac{1}{K_A} \sum_{i=K-K_A+1}^K E[f_{i:K}]} \quad (10)$$

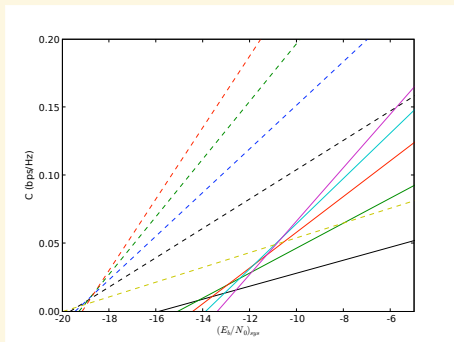
$$\mathcal{S}_0 = \frac{2\mathbf{C}'(0)^2}{-\mathbf{C}''(0)} \quad (11)$$

$$= \frac{2 \left(\int_0^\infty \frac{y}{1+\text{SNR}y} dF_y(y) \Big|_{\text{SNR} \rightarrow 0} \right)^2}{\int_0^\infty \left(\frac{y}{1+\text{SNR}y} \right)^2 dF_y(y) \Big|_{\text{SNR} \rightarrow 0}} \quad (12)$$

$$= \frac{2\mathbb{E}[y]^2}{\mathbb{E}[y^2]} \quad (13)$$

$$= \frac{2 \left(\sum_{i=K-K_A+1}^K \mathbb{E}[f_{i:K}] \right)^2}{\sum_{i=K-K_A+1}^K \left(\frac{\mathbb{E}[s^2]}{\mathbb{E}[s]^2} \mathbb{E}[f_{i:K}^2] + \sum_{j=K-K_A+1, j \neq i}^K \mathbb{E}[f_{i:K} f_{j:K}] \right)}. \quad (14)$$

Increase K_A , lose in $(E_b/N_0)_{\min}$, win in wideband slope.



For larger K you lose less! (solid $K = 10$, dashed $K = 1000$)

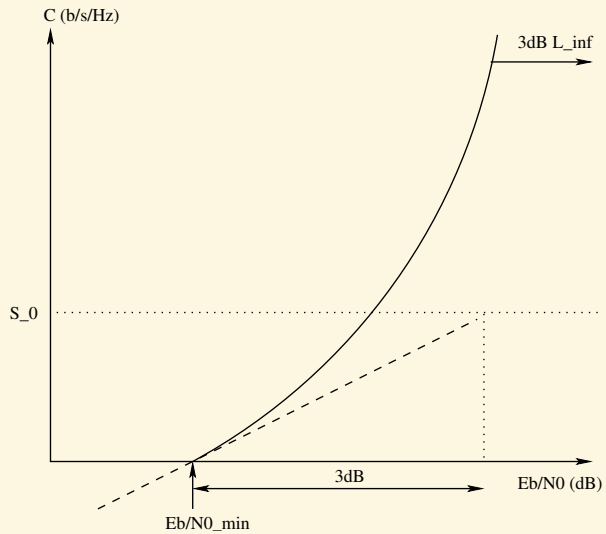
What is happening

- For larger K , the difference of K_A largest f_k diminishes
- Second moment of y inversely proportional to K_A^2
- For large K_A :

$$\frac{2\mathbb{E}\left[\frac{1}{K_A}\sum_{k\in\mathcal{A}}d_k\right]^2}{\mathbb{E}\left[\left(\frac{1}{K_A}\sum_{k\in\mathcal{A}}d_k\right)^2\right]}\rightarrow\frac{2\mathbb{E}[\mathbb{E}[d]]^2}{\mathbb{E}[\mathbb{E}[d^2]]}=2$$

Taking $K_A = \gamma K \rightarrow \infty, \gamma \leq 1$

- $(E_b/N_0)_{\min} = \frac{\log(2)}{E[s]E_{\mathcal{A}}[f]}$
- $\mathcal{S}_0 = 2$



$$\mathcal{L}_\infty = \lim_{\text{SNR} \rightarrow \infty} \left(\log_2 \text{SNR} - \int_0^\infty \log_2 (1 + y \text{SNR}) dF_y(y) \right) \quad (15)$$

$$= -\mathbf{E}_y [\log_2(y)] \quad (16)$$

$$\leq -\mathbf{E}_s \left[\log_2 \left(\frac{1}{K_A} \sum_{i=1}^{K_A} s_i \right) \right] \quad (17)$$

$$= -\mathbf{E}_{f_{K-K_A+1:K}} [\log_2(f_{K-K_A+1:K})]. \quad (18)$$

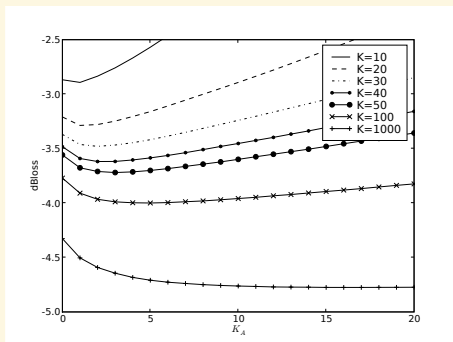


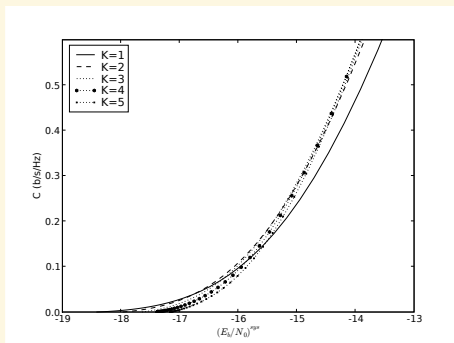
Larger K_A gives

- smaller penalty due to (17) – near-far gain (Jensen)
- larger penalty due to (18) – lost diversity gain.

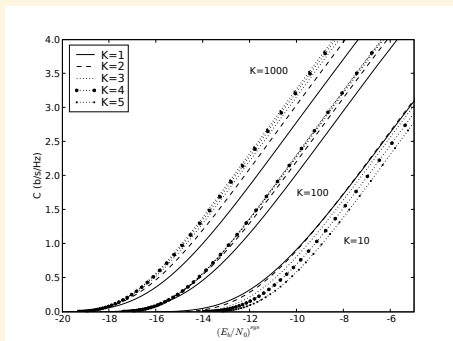
Also

- Near-far gain is bounded.
- Multiuser diversity loss is a function of user population size.
- There is a K -dependent optimal K_A .





Exact behavior at low spectral efficiency ($K = 100$).



Already with 10 users scheduling 2 at a time is ok.

- Asymptotic tools can simplify your life.
- Minimum E_b/N_0 , wideband slope and dB penalty useful parameters to analyze
- Scheduling multiple users with group-PFS
 - increases wideband slope
 - increases minimum E_b/N_0
 - decreases dB penalty with high enough K