

Cloud-native and programmable radio access networks

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Mobile Cloud Networking

The 5G Infrastructure Public Private Partnership

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Economics of mobile are changing

Softwarization and Commoditization

Software implementation of network functions on top of GPP with no or little dependency on a dedicated hardware

Full GPP vs. accelerated vs. system-on-chip

Programmable RF

Virtualization and Cloudification

- Execution of network functions on top of virtualized computing, storage, and networking resources controlled by a cloud OS.
- Share I/O resources among multiple guests

Emergence of rich ecosystem and opensource for telecom

- NFV, SDN and MEC
- Open APIs and standardized I/F



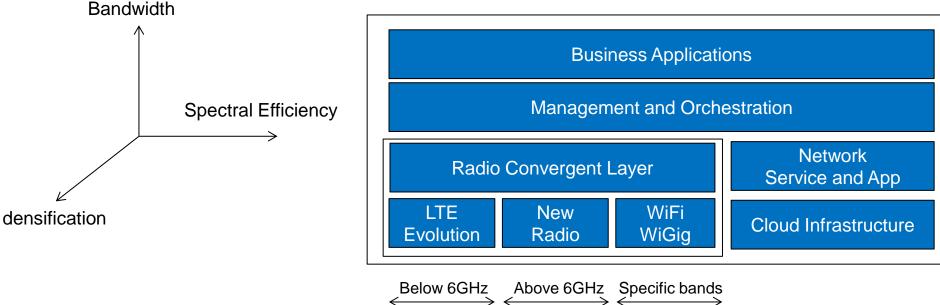


Not all of these Requirements need to be satisfied simultaneously

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5G will be a paradigm shift

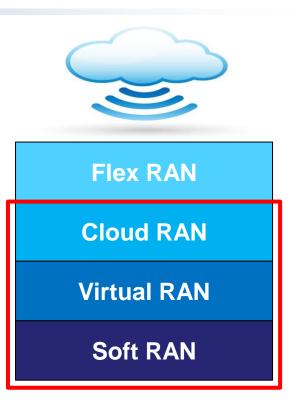
Overall 5G Solution



 5G is not just a new radio/spectrum, but also a new architecture and business helper



- Technology
- Challenges
- Results
- Conclusion





Cloud Computing

Cloud Computing disrupts IT consumption and operations

- on-demand, self-service, elastic, pay-as-you-go, metered service
- Additional advantage: automated management, remote access, multi-tenancy, rapid deployment and service provisioning, load-balancing

New business models based on sharing

Public, private, local, remote, community, and hybrid clouds

Promising business potentials (CAPEX/OPEX)

Start small and grow on-demand

Software as a service

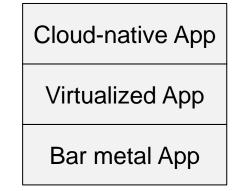
Virtual desktop, games, analytics, ...

Platform as a service

Data base, web service, ...

Infrastructure as a service

VM, storage, network, load balancer





GP-Cloud computing vs C-RAN applications

	GP-Cloud Computing	Cloud RAN
Data rate	Mbps, bursty,	Gbps, stream
Latency / Jitter	Tens of ms	< 1, jitter in ns
Lifetime of data	Long	Extremely short
Number of clients	Millions	Thousands – Millons
Scalability	High	Low
Reliability	Redundancy, load balancing	Redundancy, Offloading / load balancing
Placement	Depends on the cost and performance	Specific areas
Time scale (operation, recovery)	Non-realtime	Realtime



Cloud RAN Primer

Main idea:

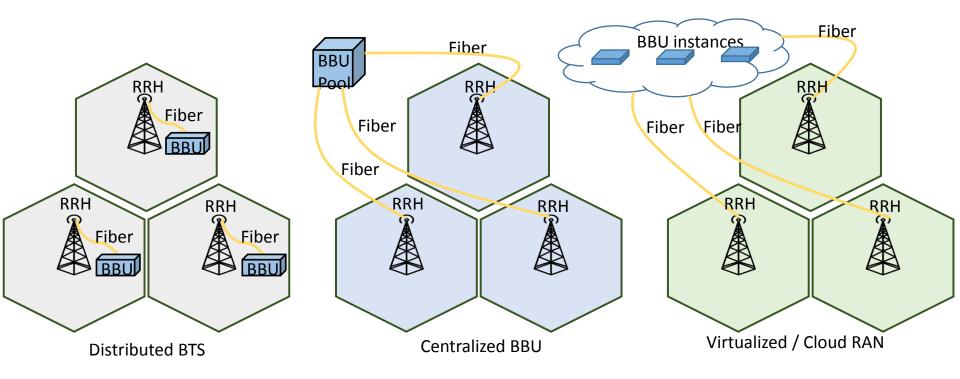
- Decouple the base station processing from the radio unit
- Perform the processing at the high performance cloud infrastructure
- > Transport the data through a high speed medium

Components

- Remote radio head (RRH): lightweight (passive) radio element with power amplifier and antennas
- Base band Unit (BBU): a centralized pool of virtualized base station covering a large set of cells (10 – 100)
- Fronthaul (FH): data distribution channel between the BBU pool and RRHs



Cloudification of RAN





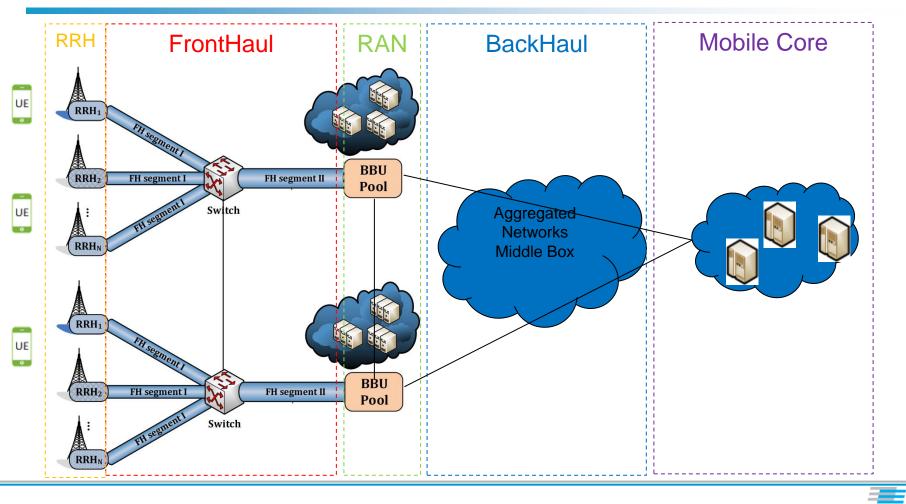
Comparison

Traditional BS, Distributed BS, and C-RAN

Architecture	Radio and BaseBand	Advantages	Drawbacks
Traditional BS	Co-located at the cell site In-BS processing	-	High power consumption Underutilized resources
Distributed BS	Split of BBU from RRH Group of RRH	Lower power consumptions Better placement of RRH	Underutilized resources
C-RAN	Split of BBU from RRH. Network of RRHs. Collocated BBUs	Even lower power consumption Better placement of BBU and RRH Lower the number of BBU Simpler network densification Rapid network deployment	Fronthaul Capacity requirement (non commodity)



Typical cloud RAN Deployment





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Benefit of a Cloudified RAN

Cooperation

- Coordinated signal processing
- Joint scheduling
- Interference management through channel feedbacks

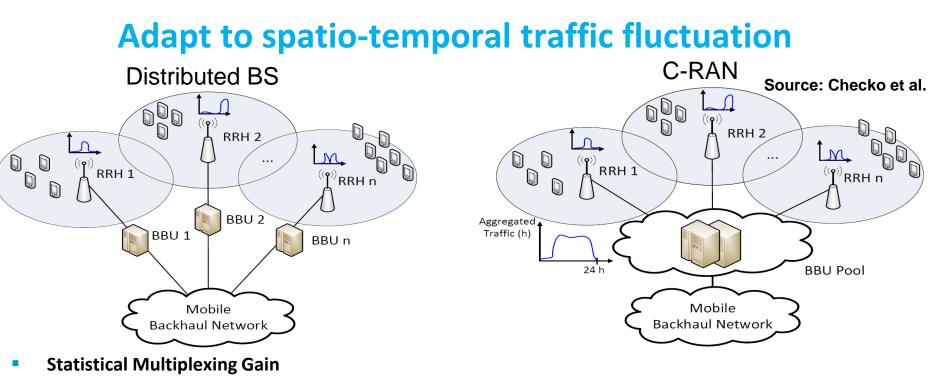
Interconnections

- Maximize statistical multiplexing gain
- Load balancing

Clustering

- RRH aggregation and assignment to BBU pools
- Reduce the number of BBUS to save energy





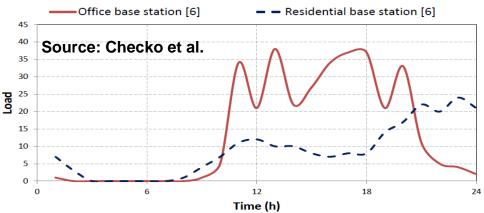
- Scalability
 - Elasticity (scale up/down)
 - Workload sharing (scale in/out)



Cloudified RAN Benefit

Exploit workload variations through the statistical multiplexing gain among multiple RAN

- BS are often dimension for the peak traffic load!
- Peak traffic load →10x off-thepeak hours
- Exception: load-aware BS



- Observation: Centralized BBUs' processing $< \Sigma$ of BSs' processing
- Statistical multiplexing gain = Σ of BSs' processing
 Centralized BBUs' processing
- Gain: depends on traffic pattern, BBU to RRH mapping, BBU load balancing



Cloudified RAN Benefit

Improve of spectral efficiency (throughput, latency)

 Centralization of BBU pool in C-RAN facilitates the inter BBU cooperation

Joint scheduling

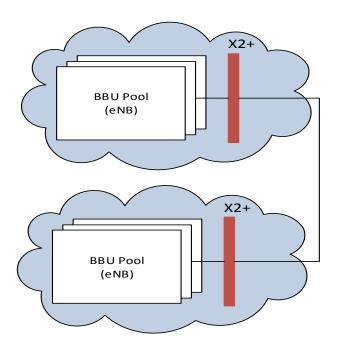
Inter cell interference (e.g. elCIC)

Joint and coordinated signal processing

utilize interference paths constructively
 (e.g. CoMP, MU-MIMO)

Shared Context

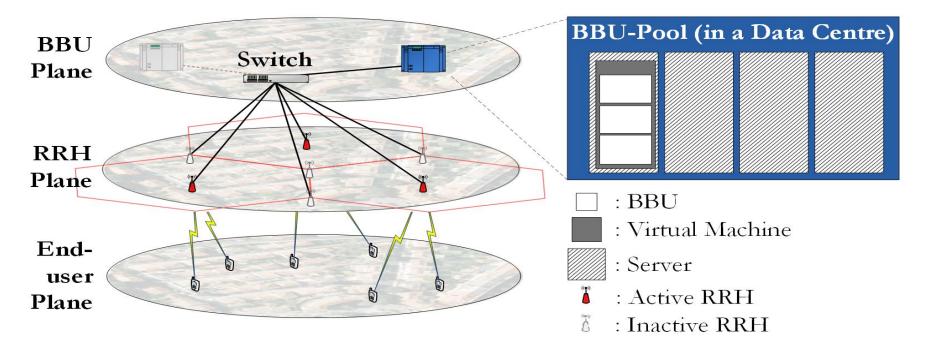
reduce control plane signaling delay (e.g. handover, co-scheduling via X2+)





Cloud-RAN Example

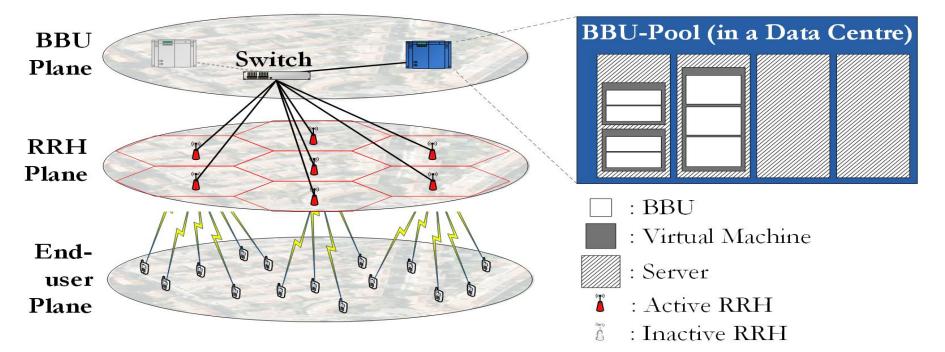
 With few users, 3 RRH-BBU pairs cover the service area and provide the requested capacity.





Cloud-RAN Example

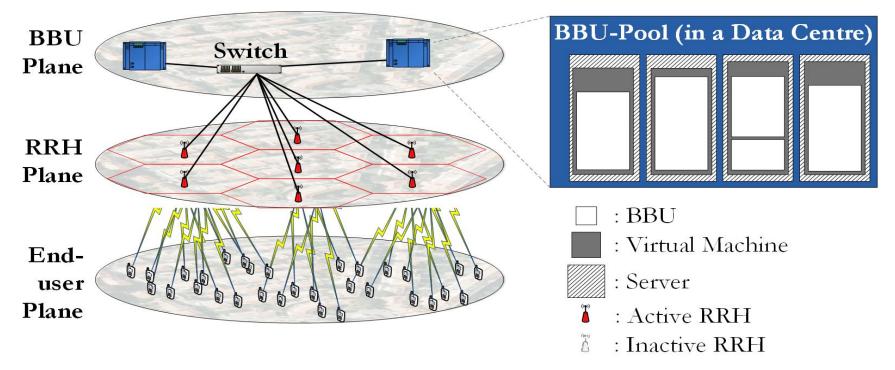
 With more users, extra RRHs are activated and BBUs instantiated, to provide the requested capacity.





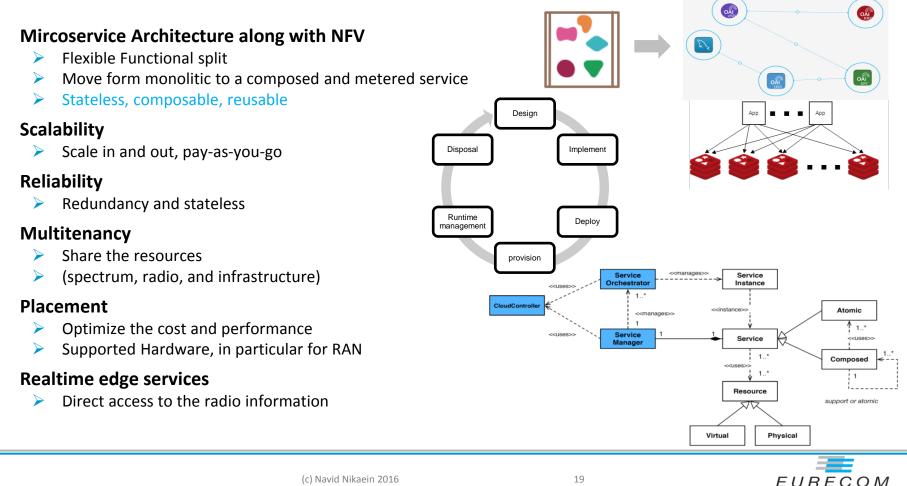
Cloud-RAN Example

 With more users, extra RRHs are activated and BBUs instantiated, to provide the requested capacity.

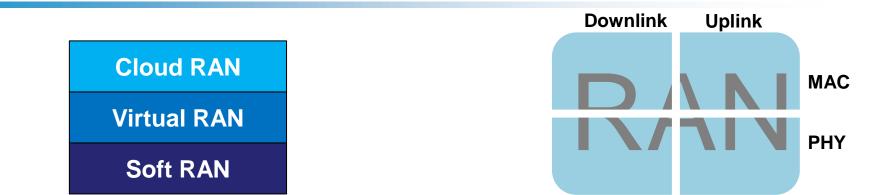




Cloud-Native RAN



Cloud-RAN Challenges



- Capacity, latency, and jitter requirements for fronthaul
- BBU processing budget and protocol deadlines
- Realtime, virtualization environment and BBU performance
- Active RRH and Flexible Functions Split
- E2E Service modelling and template definition
- NFV Service manager and orchestrator



Transport Network between RRH and BBU

- Dark fiber
- WDM/OTN: Wavelength-division multiplexing (WDM)/Optical Transport Network (OTN)
- Unified Fixed and Mobile access (microwave)
- Carrier Ethernet

Protocols

- Common Public Radio Interface (CPRI)
- Open Base Station Architecture Initiative (OBSAI)
- Open Radio equipment Interface (ETSI-ORI)

Key requirements

Supported Topology (star, ring, mesh), reliability, distance, multiplexing, capacity, scalability



Latency required by the HARQ RRT deadline

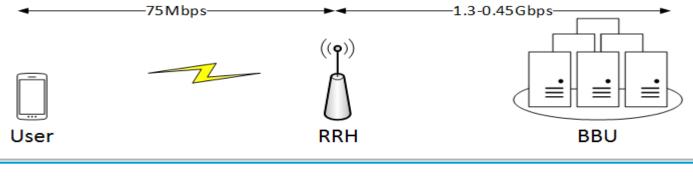
250 us maximum one-way latency adopted by NGMN, limiting the length of BBU-RRH within 20-40 Km (given that the speed of light in fiber is approximately 200x10⁶ m/s)

Jitter required by advanced CoMP schemes

- <65 ns(MIMO, 36.104) timing accuracy in collaboration between base stations, which is the tightest constraint.
- Frequency error < 50 ppb (macro BS)</p>
- ➢ BER < 10e-12</p>

20MHz channel BW, SISO, 75 Mbps for users

2.6Gbps on Fronthaul without compression (0,87Gbps with 1/3)





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 $C = 2 \cdot N_{Antenna} \cdot M_{Sector} \cdot F_{Sampling} \cdot W_{I/Q} \cdot C_{carriers} \cdot O_{coding+proto} \cdot K_{comp}$

Bandwidth	Nantenna	M_{sector}	$F_{sampling}$	$W_{I/Q}$	$O_{coding+proto}$	$C_{Carriers}$	K	Data Rate
1.4MHz	1x1	1	1.92	32	1.33	1	1	163Mb/s
5MHz	1x1	1	7.68	32	1.33	1	1	650Mb/s
5MHz	2x2	1	7.68	32	1.33	1	1	1.3Mb/s
10MHz	4x4	1	15.36	32	1.33	1	1/2	2.6Gb/s
20MHz	1x1	1	30.72	32	1.33	1	1	2.6Gb/s
20MHz	4x4	1	30.72	32	1.33	1	1/3	3.4Gb/s
20MHz	4x4	1	30.72	32	1.33	1	1	10.4Gb/s

Costs

➤ Tens of BS over long distance → 100 Gbps

Savings

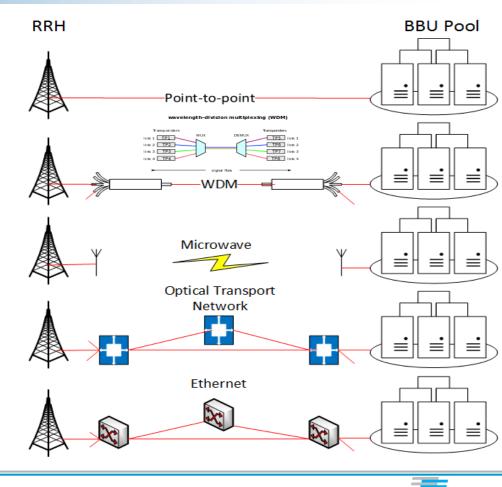
- Equipment's
- Energy



Medium	Bit rate	Distance	Remark
Fiber	100Gbps	~20Km	OTN: expensive
Copper	10Gpbs	100m	Low cost, SYNC
Wireless	1Gbps	2-15Km	LoS, high latency

Synchronization

- Frequency of transmission
- Handover, coding
- Solution
 - > GPS
 - PHY layer clock, SyncEth
 - Packet-based sync (IEEE 1588v2)



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Asynchronous Ethernet

- Reduce the fronthaul capacity
- I/Q transport over Ethernet
- Some DSP in RRH to reduce transport speed/cost (split)
 - Decoupling of user-processing and cell-processing (iFFT/FFT)

Advantages

- Cost saving (reuse, commodity hardware)
- Switching (packet-based)
- Multiplexing / load balancing
- Flexible topology (mesh)

Challenges

- Distributed computation
- Cheap synchronization ((GPS, 1588v2)
- Real-time I/Q over Eth links (copper, low-cost fiber)

RRH **BBU Pool** \equiv ≣ Atomic Clock Small-cel RRH Small-cel L1L RRH 40 GbE (switched) L1L 1GbE (IF3/4) Small-cel RRH 10GbE (IF3) CPRImacro-cel CPR RRH L1L

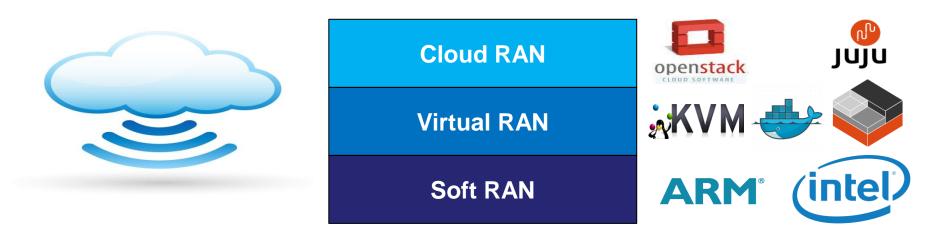
Hot topics

- > IEEE 1904.3 encapsulation and mapping of IQ data over Ethernet
- IEEE 802.1 CPRI fronthaul discussion with Time Sensitive Networking task force
- > CPRI → CPRI2?
- 3GPP proposal on a study item on variable rate multi-point to multi-point packet-based fronthaul interface supporting load balancing



Soft RAN BBU processing budget

- 4G Feasible on General Purpose Processors (x86)
- An eNB is approximately 1-2 x86 cores on Gen 3 Xeon silicon
 - > Perhaps more power efficient solutions from TI, Freescale or Qualcomm
 - But: lose commodity software environment and common HW platform to high-layer protocols and cloud





Soft RAN



BBU processing budget for peak rate

eNB Rx stats (1subframe)

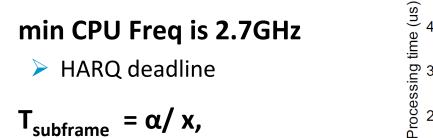
- OFDM demod : 109.695927 us
- ULSCH demod: 198.603526 us
- ULSCH Decoding : 624.602407 us

eNB Tx stats (1 subframe)

- OFDM mod : 108.308182 us
- DLSCH mod : 176.487999 us
- DLSCH scrambling : 123.744984 us
- DLSCH encoding : 323.395231 us
- → 730 us (< 1core)</p>

- → 931 us (<1 core)</p>
- Efficient base band unit is challenging
- With AVX2 (256-bit SIMD), turbo decoding and FFT processing will be exactly twice as fast
 - <1 core per eNB</p>
- Configuration
 - gcc 4.7.3, x86-64 (3 GHz Xeon E5-2690),
 - 20 MHz bandwidth (UL mcs16 16QAM, DL mcs 27 64QAM, transmission mode 1 SISO)
 - 1000 frames, AWGN channel





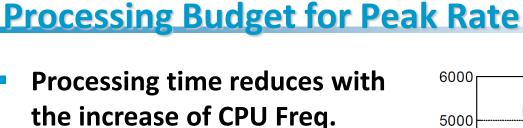
 $\geq \alpha = 8000$

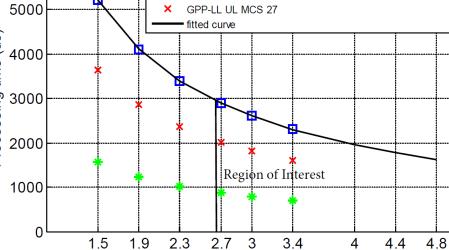
Soft-RAN

x is the CPU freq GHZ

min CPU Freq is 2.7GHz

HARQ deadline





CPU Frequency (GHz)

Total RX/Tx processing vs. CPU frequency

GPP-LL DL MCS 27

Note: FDD LTE HARQ requires a round trip time (RTT) of 8ms

 \succ Tx+RX \leq Tharq/2–(acquisation+transport+offset) \approx 3ms

~2ms RX and 1ms TX (can't be fully parallelized)



inside

Xeon

RFACE

GPP-LL Total Tx/Rx processing(DL MCS 27, UL MCS 27)

Soft RAN Considerations

Key Consideration to meet the deadlines (SF, protocol)

- Real-time OS (linux with deadline scheduler) and optimized BIOS
 - Problem: OS scheduler latency (kernel is not pre-emptible)
- Real-time data acquisition to PC
- SIMD optimized integer DSP (SSE4, AVX2)
- Parallelism (SMP)
- ≻ x86-64
 - more efficient for Turbo decoding because of the number of available registers is doubled

Remove bottlenecks with

- hardware accelerators or hybrid CPUs
 - Turbo decoders (easily offloaded to FPGA-based accelerators), FFT, PDCP (de)enryption
- GPUs or Xeon PHY-type devices
 - Perhaps interesting for Turbo-decoders and encoders than FFT
 - Main issue in both FPGA/GPU offloading
 - High-speed low-latency bus between CPU memory and external processing units



Realtime, virtualization environment and BBU performance RTOS issues

- Low-latency radio applications for PHY (e.g. 802.11x,LTE) should run under an RTOS
 - Meet strict hard deadline to maintain the frame/subframe and protocol timing
 - efficient/elastic computational resources (e.g. CPU, memory, network)

Example OS

- eCos/MutexH for generic GNU environment
- RTAI for x86
- VXWorks (\$\$\$)
- Example: RTAI / RT-PREEMPT kernel can achieve worst-case latencies below 30µs on a loaded-PC. More than good enough for LTE, but not for 802.11x because of MAC timing.
- Should make use of POSIX multithreading for SMP
 - Rich open-source tool chains for such environments (Linux, BSD, etc.)
 - Simple to simulate on GNU-based systems for validation in user-space
 - Allow each radio instance to use multiple threads on common HW



Realtime, virtualization environment and BBU performance Issues with standard Linux Kernels

Scheduler latency

- > Kernel is not pre-emptible
- Overhead in disabling/enabling interrupts
- Mainstream kernel solutions, the RT-Preempt patch and out-ofthe-box Linux kernel (>3.14) converts Linux into a fully preemptible kernel
 - Kernel preemption (RT-PREEMPT) mainstream until 2.6.32 (patches afterwards)
 - Latency reduction (soft-RT kernels) with DEADLINE_SCHED Version >3.14

Patches / dual-OS solution

ADEOS + RTAI/Xenomai



Realtime, Virtualization environment and BBU performance

- Virtual Machine (VM) e.g., KVM: XVM
 - > A complete OS is deployed as a guest
 - Virtualisation layer that emulates physical resources
 - Hypervisor that manages requests for CPU, memory, hard disk, network and other hardware resources

Virtualisation Environment (VE) – e.g., LXC and Docker:

> No hardware emulation nor hypervisor and guest OS (containers).

- Use and share the OS and potentially device drivers of the host
- OS scheduler manages the request for physical resources



Realtime, Virtualization environment and BBU performance

General Purpose Platform (GPP)

dedicated machine.

Kernel-based Virtual Machine (KVM)

Linux virtualisation infrastructure that turns it into a hypervisor.

Linux Container (LXC)

operating-system-level capability for running isolated Linux Virtual Environments (VE) on a single control host.

Docker

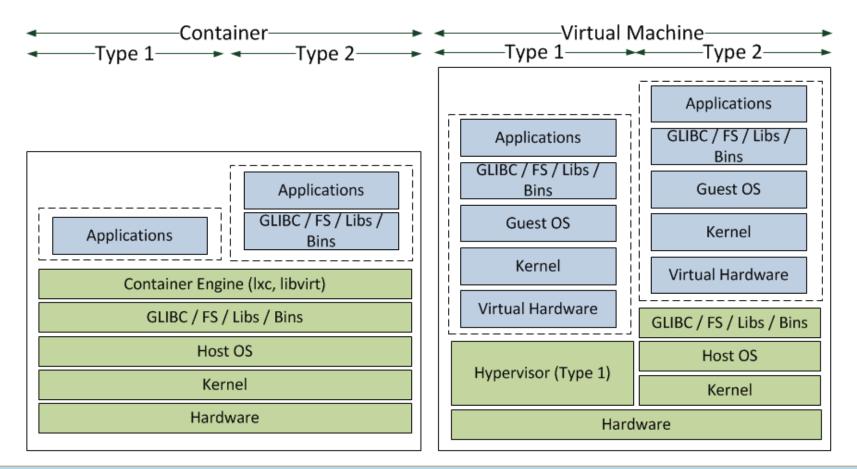
LXC-based portable container engine that encapsulates an application with all its dependencies.

Options:

Iow latency kernel, prioritization of processes.



Realtime, Virtualization environment and BBU performance

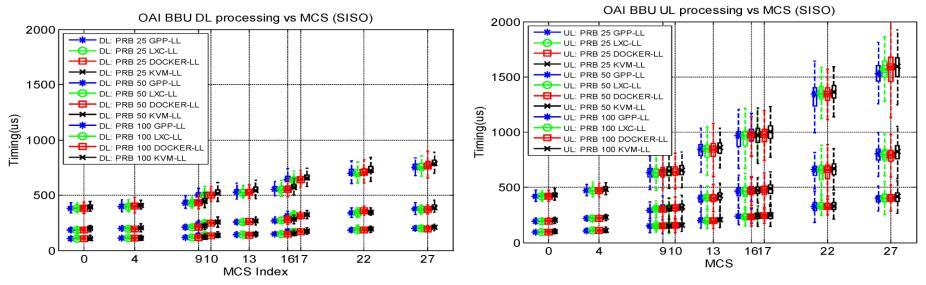




Virtual-RAN Processing Budget for Peak Rate KVM

DL and UL BBU processing load for various MCS, PRB, and virtualization flavor

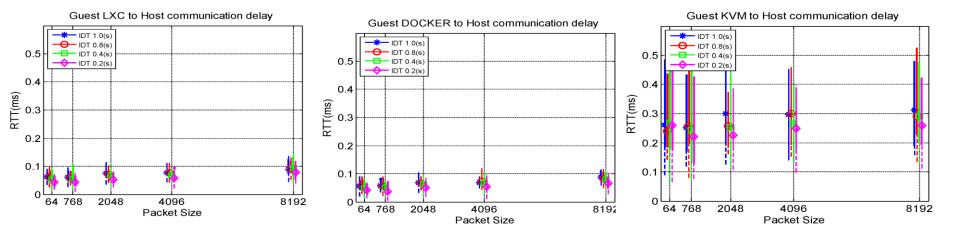
Comparable BBU Processing time





Virtual-RAN Additional Consideration

- I/O access delay
 - RF, ETH, and HW accelerator
 - RF Passthrough vs Hardware virtualization (and sharing)
 - Delay and jitter requirement on the fronthaul network
- Limitation of the guest-only network data rate





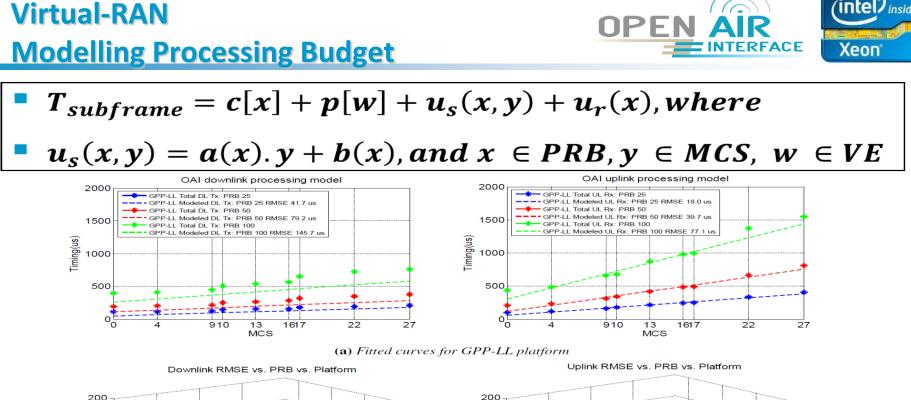
inside

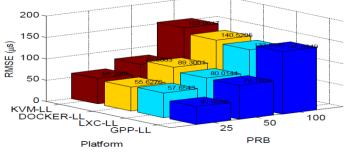
Xeon

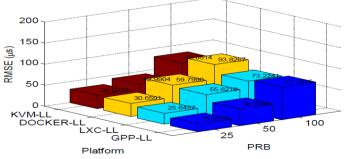
FACE

OPEN

×K









nside

Cloud RAN CPRI-based RRH

CPRI is

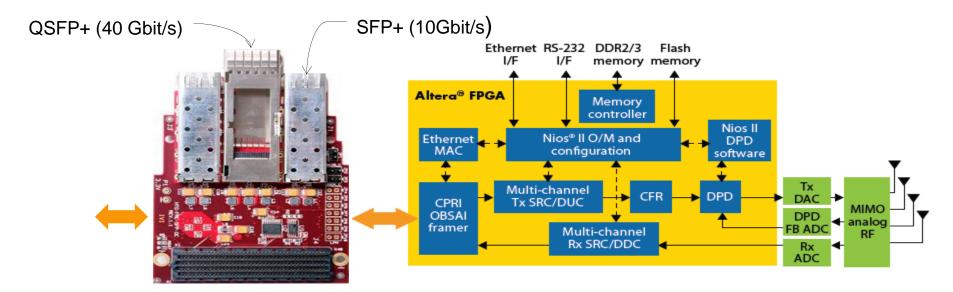
- A synchronous protocol for high-speed transport of I/Q baseband signals between BBU and RRH
 - Uses Gigabit ethernet-like (10,40,later 100) physical links based on 122.88 MHz clock and optical transport (for 40,100)
 - \checkmark Line rates up to 9.8 Gbit/s (20 MHz antenna port \approx 1.2 Gbit/s bi-directional)
 - All RRH are driven by common clock from BBU => tight synchronization in time/frequency is possible
 - Framing is scalable to allow for different number of antennas and channel bandwidths
- I/Q transfer is standardized and flexible (number of bits, sampling rate, etc.)
- RF control allows for proprietary signaling to control RF (biggest issue for developers in order to adapt to different RRH vendors)



Cloud RAN CPRI-based RRH

 CPRI-based RRH are usually built using FPGA (Xilinx/Altera) platform with small embedded system

Coupled with RF cleanup (upsampling/downsampling filters, TX predistortion)





Cloud RAN

Active RRH and Ethernet Frontahul

CPRI-gateways (switches)

- One end is Ethernet (connection with BBU-pool) other is CPRI for commercial RRH
- Possibility to use a CPRI-GW to deliver synchronous I/Q to group of RRH (P2MP or multi-hop) from a common Atomic reference and provide generic Ethernet to BBU-pool

"Cheap" RRH (e.g. large indoor networks)

- Regular Ethernet or (syncE) +1588v2 (even copper!)
- Low-power (<20W), cheap I/Q transport to BBU (i.e. not CPRI) with copper or cheap-fiber Ethernet</p>
- Some DSP in RRH to reduce transport speed/cost
- Low-cost RF (e.g. Existing Lime microsystems-based PCIe solution)
- Open architecture synchronization solution

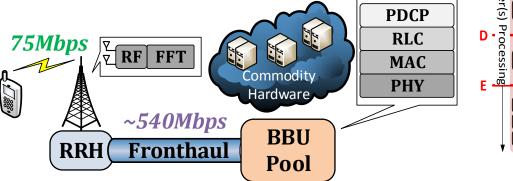
BBU is slave to network of RRH

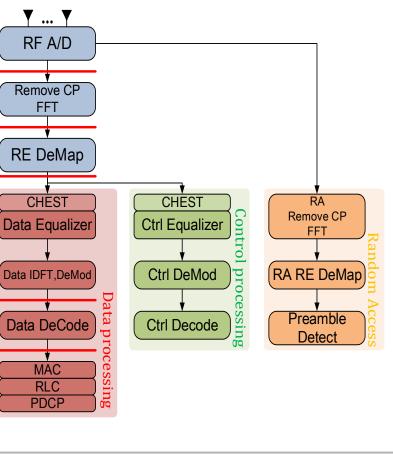
Cloud RAN

Active RRH and Flexible Functions Split

Place more BBU processing at the edge of the network

- Reduce FH capacity requirement
- Add FFT and remove CP at RRH almost halves the FH bandwidth
 - From 1Gbps to 540Mbps
- However, some disadvantage...
 - Expensive RRHs
 - Less coordination







A • @

Processing

B

C –

User(s)

Process

FFT

Cloud RAN <u>Active RRH and Flexible Functions Split</u>

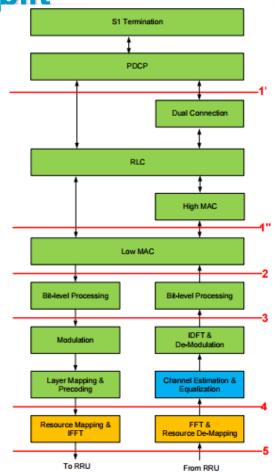
- China Mobile/NGFI approach
- Similar to small-cell forum

	Interface 1		Interface	e 2	Interface 3 Ir		Interface	Interface 4 Interface 5		5
	Bandwidth	Ratio	Bandwidth	Ratio	Bandwidth	Ratio	Bandwidth	Ratio	Bandwidth	Ratio
Downlink	174 Mb/s	1	179.2 Mb/s	1	125.2 Mb/s	1	498 Mb/s	3	9,830.4 MB/s	66
Uplink	99 Mb/s	1	78.6 Mb/s	1	464.6 Mb/s	6	2,689.2 Mb/s	36	9,830.4 MB/s	131

Table 3-1: Maximum Interface Bandwidth

Table 3-2: Interface Delay

Interface 1		Interface	2	Interface	3	Interface	4	Interface	5
Delay	Ratio	Delay	Ratio	Delay	Ratio	Delay	Ratio	Delay	Ratio
Less than 100 ms	1	Less than 1 ms	100						





Cloud RAN

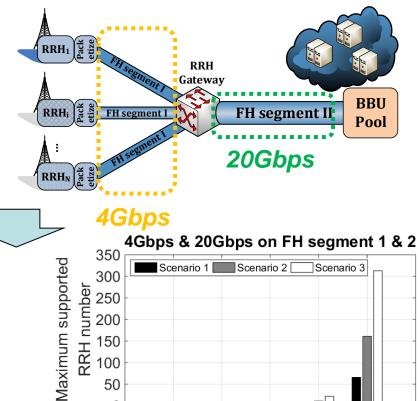
Where to split?

Derive maximum supported RRHs based on achievable peak-rate

Scenario	1	2	3			
Bandwidth	20 MHz					
Oversampling Ratio	1					
Rx Antennas	4					
Cyclic prefix length		Normal				
MIMO	4 Layer					
PUCCH RB	4					
SRS BW Config	7					
SRS SF Config	9					
Control Overhead		4.3%				
RA Config	0					
RA Overhead		0.3%				
Modulation	64 QAM	16 QAM	QPSK			
TBS index	26	16	9			
Time sample bitwidth	16					
Frequency sample bitwidth	16					
LLR bitwidth	8					



Scenario	1	2	3		
Split A	5				
Split B	8				
Split C		9			
Split D	7	11	22		
Split E	66	161	313		





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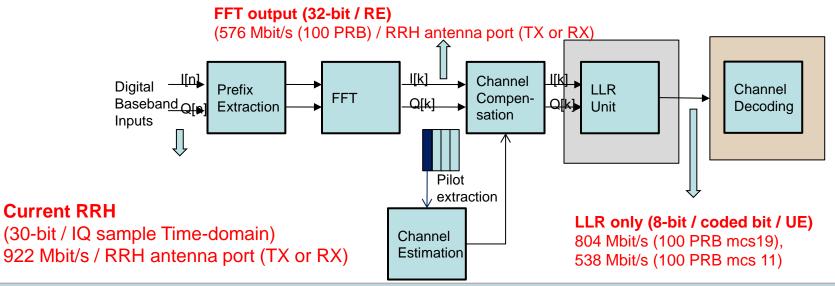
Cloud RAN

Where to split?

- TX
 - Full L1 TX in RRH
 - MAC (scheduler) must provide
 - Transport channel SDUs (common and dedicated)
 - Any precoding information for TM7-10

RX split is still under investigation

- Depends on number of UEs / RE / RRH (i.e. MU detection per RE)
- > And on models for realistic uplink resources (average MCS) in dense deployments





Cloud-RAN Where to split?

It is not clearly evident that transport of the quantized LLR provides significant savings

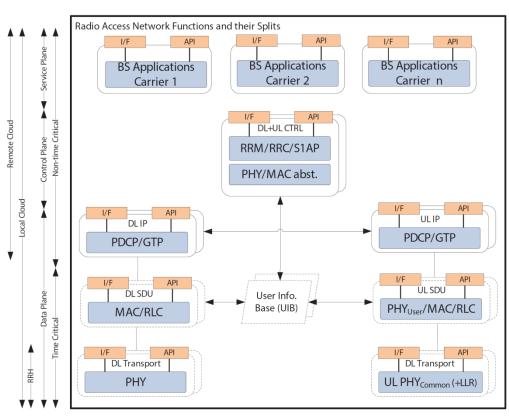
However

- > The assumption here is that no further compression is required
 - Because of the quasi-discrete nature of the LLRs, further compression could bring savings
- If compression can bring us below 8-bit/coded bit/UE then
- Also, we can trade-off some performance by quantizing LLRs to 4bits, then there would be significant fronthaul savings
- TX fronthaul rates can be significantly reduced if baseband TX is performed in RRH
 - Could be interesting for densely-deployed DL-only RRH



Cloud-RAN Where to split?

- RRC and MME Placement
- PDCP as a convergent layer
- PHY_{user} as a variable
 W and W/O MAC/RLC
- Allow split across RRH, local, and remote cloud
 - I/F ➤ Orchestration logic
- - > Controller logic

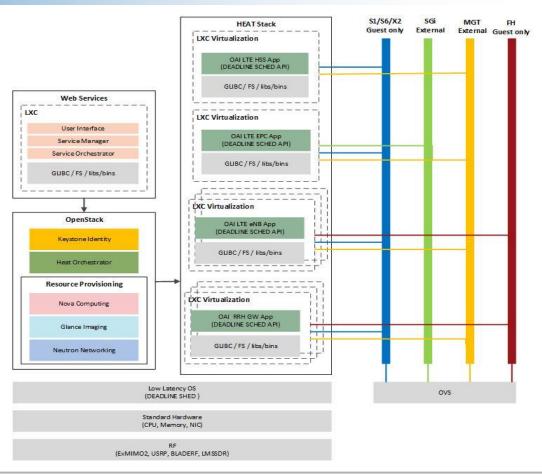




C-RAN Testbed on Sophia Antipolis Campus

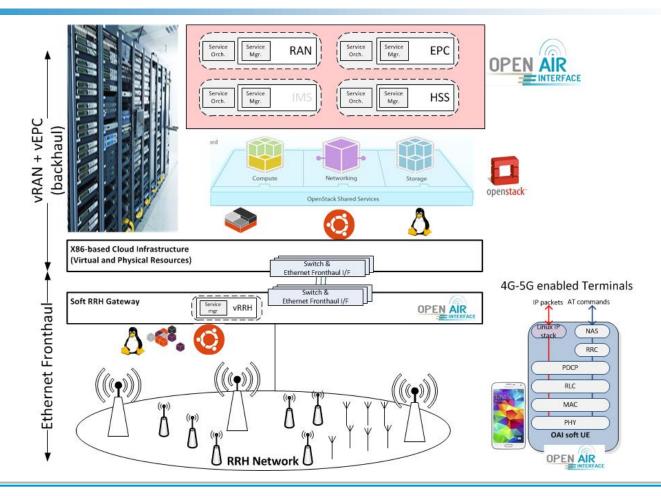
Three components

- web service
- > OpenStack
- Heat stack
- Heat Template describes the virtual network deployment
- Linux Container
- Open vSwitch
- Low latency kernel
- RF frontend HW





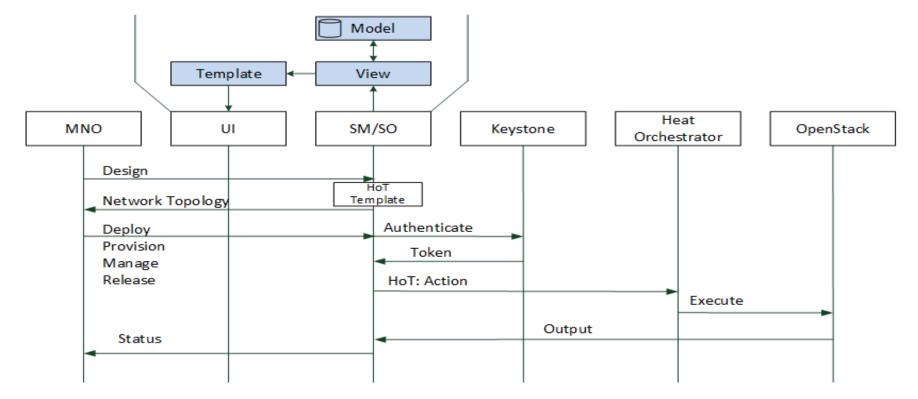
C-RAN Testbed on Sophia Antipolis Campus





C-RAN Testbed on Sophia Antipolis Campus <u>Message Sequence (Openstack)</u>

Orchestrator is key in the life cycle management





C-RAN Testbed on Sophia Antipolis Campus Heat Orchestration Template (HOT)

- The instantiation of a whole system (e.g., an LTE ecosystem) can be easily achieved with HOT
 - virtual components of the communication network defining a network slice
- Different level of abstractions are required





C-RAN Testbed on Sophia Antipolis Campus Heat Orchestration Template (HOT) - Example

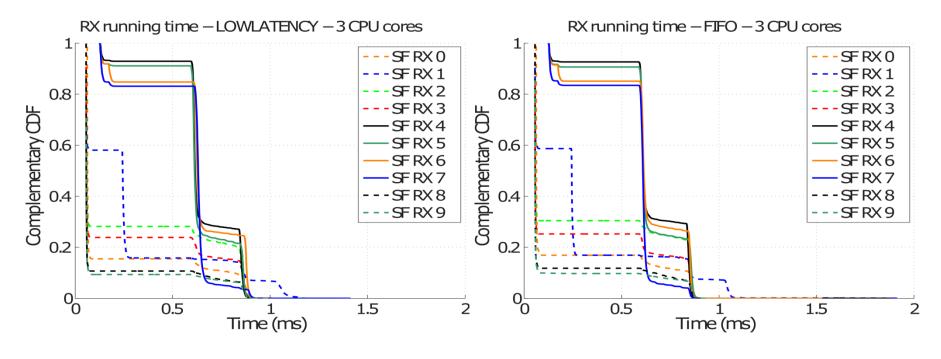
```
description: LTEaaS,
parameters: { key_name: {
               type: string, description: Name of a KeyPair to enable SSH access to the instance, default : cloud,
                                                                                                                            }},
               HSS: {... } },
resources: {
               EPC: { ... } },
               ENB: {
                      type: OS::Nova::Server,
                      properties: {
                               image: enb-1,
                               flavor: eNB.med,
                               key name: cloud,
                               networks: [{network: S1, }],
                               user data:{
                                              #!/bin/bash\n
                                              MY_IP_S1=`ip addr show dev eth0 | awk -F'[ /]*' '/inet /{print $3}'`\n
                                              sed -i 's/MY IP/'$MY IP S1'/g' /etc/hosts\n
                                              sed -i 's/EPC_IP/'$EPC_IP'/g' /etc/hosts\n
                                              sed -i 's#MY_IP#'$MY_IP_S1'/24#g' enb.band7.exmimo2.lxc.conf\n
                                              sed -i 's#EPC IP#'$EPC IP'#g' enb.band7.exmimo2.lxc.conf\n
                                              build_oai.bash -I ENB -t SOFTMODEM -D --run -C enb.band7.exmimo2.lxc.conf > /tmp/oai.log\n,
                               params: {
                                              $EPC IP: {get attr: [EPC, first address],}
               }
```



}}

C-RAN Testbed on Sophia Antipolis Campus Impact of the OS scheduler

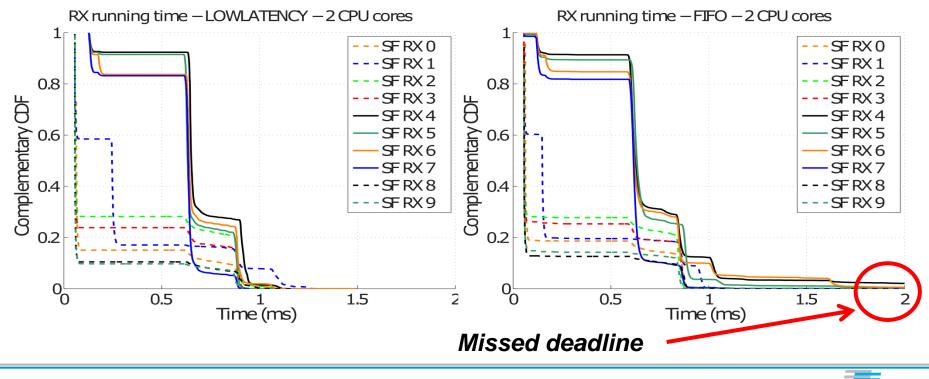
- FDD, 10MHZ, SISO, with EXMIMO RF
- UL Processing time: Only 4 uplink sub-frames
 - SF #0, 1, 2 and 3, allowing UL transmission to occur in SF # 4, 5, 6, 7.
 - Full uplink traffic





C-RAN Testbed on Sophia Antipolis Campus Impact of the OS scheduler

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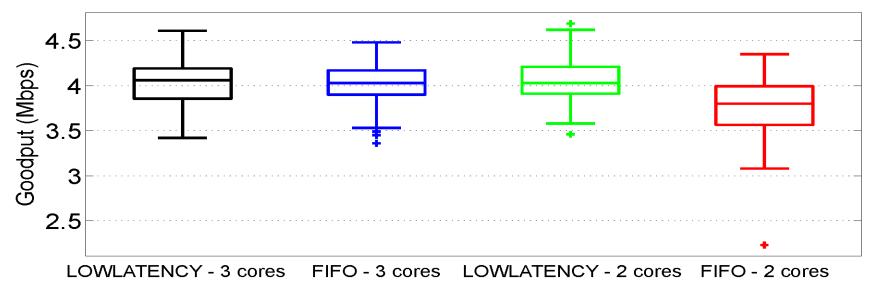




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C-RAN Testbed on Sophia Antipolis Campus Impact of the OS scheduler

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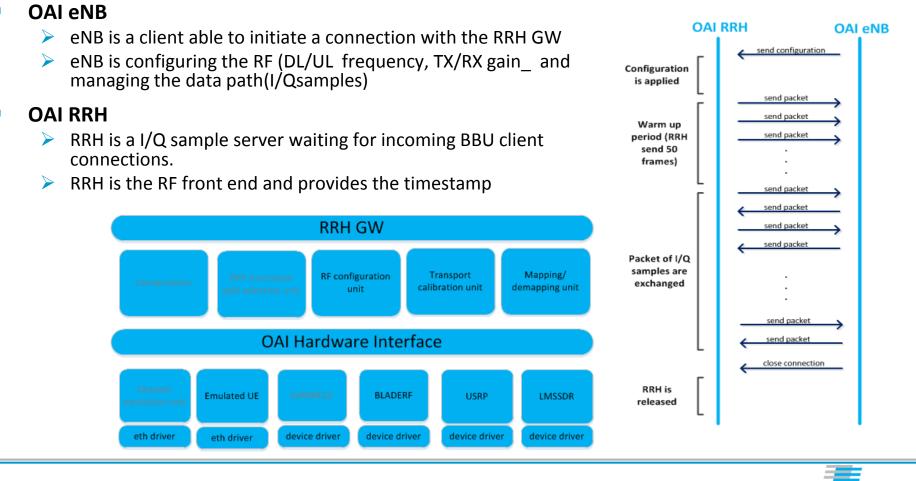


UL Channel Goodput



C-RAN Testbed on Sophia Antipolis Campus

OAL RRH



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C-RAN Testbed on Sophia Antipolis Campus OAL RRH

The FH interface is divided logically into two streams:

- > Data: transports payload, packet length is a function of BW and MTU.
- Control: in-band or out-of-band; eNB configures and manages RRHs

Two flavors of FH protocol are supported:

- UDP transport protocol: offers statistical multiplexing (multiple simultaneous communication on the same medium) at the cost of one additional layer in the protocol stack.
- RAW Ethernet: offers minimal protocol stack but unable to support statistical multiplexing

Header format (no split case)

Fields	Size (bits)	Description
Control Flag	1	This flag is used to specify whether the payload of the packet is either data or control data. value=0 ->data value=1 ->control data
Timestamp	64	The timestamp of the packet is the time that the payload was generated by the radio equipment.
Antenna ID	16	Antenna ID is a number used to map a packet to the appropriate antenna of the radio front-end equipment.



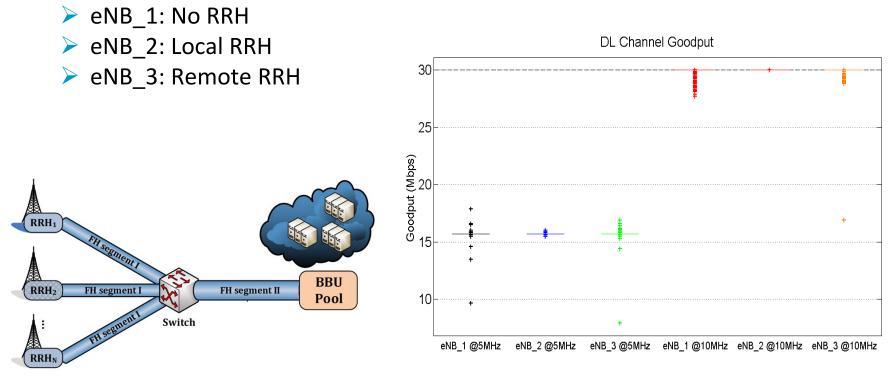
C-RAN Testbed on Sophia Antipolis Cam



intel

Xeon

Three setting (FDD, SISO, with USRP B210 RF, Eth fronthaul network)





Deployed CRAN NFV Service Template Juju

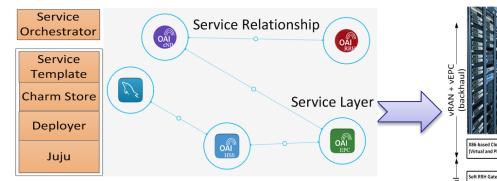


MWC 2016

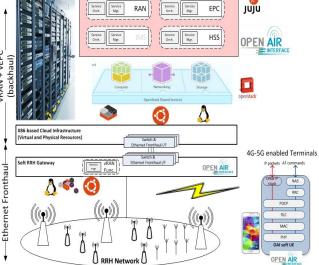


MCN/Mobicom Demo





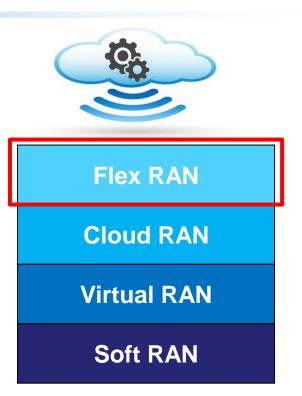
- Demo @ MWC 2016 w/ Canonical
- https://insights.ubuntu.com/2016/02/22/canoni cals-vnf-pil-for-nfv-scale-out-architectures/
- https://jujucharms.com/q/oai



Life Cycle KPI	Unit	KPI measurements
Installation	Time(s)	600 seconds
Configuration	Time(s)	4 seconds
Disposal	Time(s)	< 1 seconds
Service upgrade duration	Time(s)	122-300 seconds



- Technology
- Challenges
- Results
- Conclusion

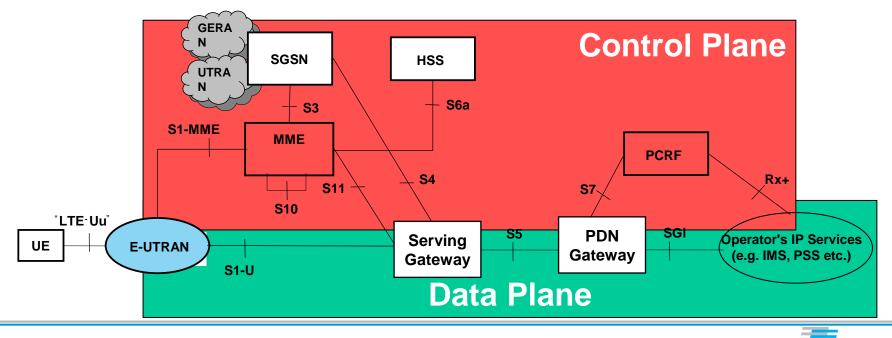




Software defined networking

Simplify network control and coordination

- > Separation of the control from the data plane with a well-defined API
- Consolidation of the control plane
- Network abstraction and programmability



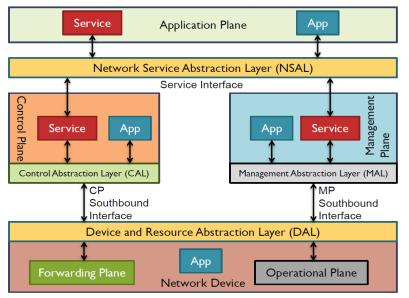
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SDN for Wireless

Standard flow-level abstraction is not enough

- Stochastic nature of wireless links
- Resource allocation granularity and time-scale
- Heterogeneity of RAN

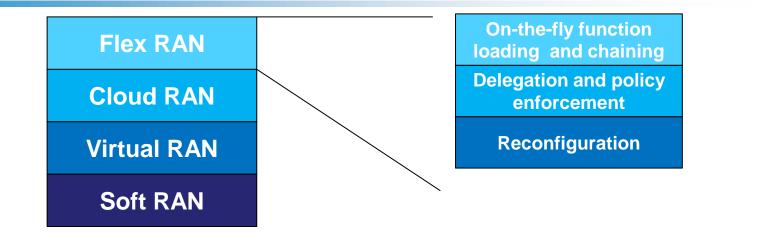
IEEE Architecture (RFC7426)



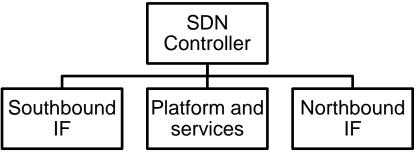
- Requirements for wireless network abstraction
 - State management
 - Resource allocation
 - Network monitoring
 - Network control
 - Network Application
 - Semantics
 - Programming SDK
 - Northbound Interface
 - Network Operating System
 - Network Abstraction
 - Southbound Interface
 - Network Infrastructure



SDN Challenges



- Radio and core API and Southbound Protocol
- Network Abstraction and graphs
- Scalability and Control delegation mechanisms
- Realtime control
- Low latency edge packet services
- Cognitive management, self-adaptive, and learning methods
- Northbound Application programming interface





Radio and core API and Southbound Protocol

- Control plane APIs allowing fine grain radio and core control and monitoring
- Platform- neutral and extendable protocol message service
 - Language agnostic
- Optimize message footprint
 - Aggregation
 - (de)serialization
- Asynchronous control channel
 - Queue
 - Pubsub communication model
- Supported network topologies
 - > P2P, P2MP, and possibly others

Message	Field	Usage
	Configuration type	Type of configuration, either set or get
	Cell configuration flag	Bit map of the requested cell configuration
Configuration	Cell configuration list	List of cells (in IDs) to request configuration
request	UE configuration flag	Bit map of the requested UE configuration
	UE configuration list	List of UEs (in IDs) to request configuration
Configuration	Cell configuration	Requested cell configuration report
reply	UE configuration	Requested UE configuration report
	Status type	Can be periodical, one-shot, event-driven
	Status period	Period in Transmission Time Interval (TTI)
Status	Cell status flag	Bit map for the requested cell status
request	Cell list	List of cells (in IDs) to request the status
	UE status flag	Bit map for the requested UE status
	UE status list	List of UEs (in IDs) to request the status
Status	Cell status	List of cell including the statistic reports
reply	UE status	List of UE including the statistic reports



Network Abstraction and graph

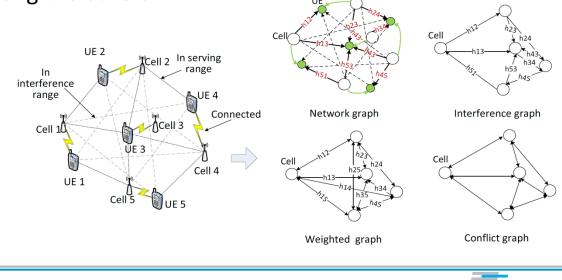
- Effective representation of the network state at different network levels allowing
 - fine-grained programmability, coordination and management of atomic or composed services across different domains/regions via

Network graphs can be separated based on

Time, Region, carrier, cell among the others

Encompass data models

- Time-frequency status and resources
- Spatial capabilities
- Key performance indicators



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Scalability, Control delegation mechanisms, and <u>Realtime Control</u>

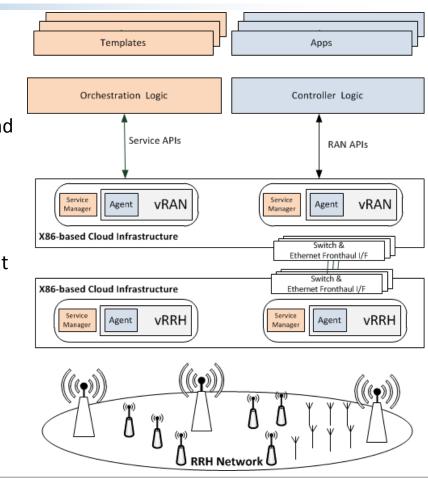
- Feasible to achieve a realtime RAN programmability at TTI level (1ms)
 - Realtime control: Guarantee a (quasi-) deterministic reaction time of a control command triggered by the controller

Hierarchical controller logic

- \succ non-time critical \rightarrow centralized entity
- \succ time critical \rightarrow edge entity
- May offloaded time critical operation to an agent acting as a local controller

Network applications

- Proactive based on periodical event
 - Scheduler
- Reactive based on event-triggering
 - E.g. mobility manager
- Interplay with the orchestrator



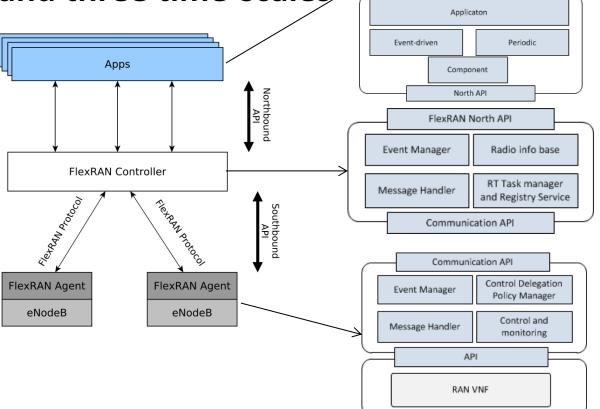


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Programmable RAN Controler-Agent Design

Three subsystems and three time-scales,

Message	Field	Usage
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Programmable RAN DL performance comparison

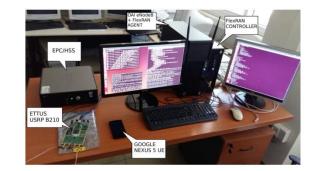
Flexan Polocor

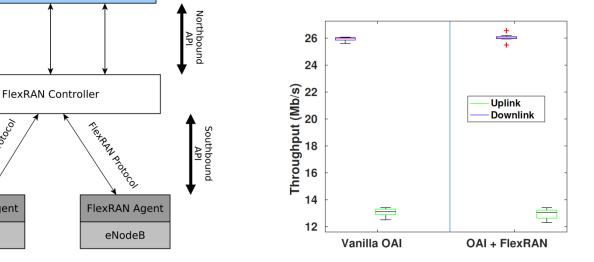
FlexRAN Agent

eNodeB

Considered controller apps

- DL scheduler
- Monitoring and analytics







Apps





Conclusion

3G

-

5G

G

Conclusion

 4G/4G+ feasible on General Purpose Processors (x86) and Virtualization environment

Exploit hybrid CPUs

Gap between virtualization and cloudification

Exploit the microservice and NFV principles

- Realtime network programmability is feasible at TTI level
 Exploit MEC principles for the data-plane programmability
- Gat between static and cognitive control and management, self-adaptive, and learning methods

> Exploit machine learning and data mining techniques



Future Research Topic

Functional split in RAN and CORE

- What is the optimal split under capacity-limited fronthaul/backhaul and processing-limited compute resources ?
- How to change the functional split on the fly?

Cognitive self network management

- What are the right network abstraction and modelling?
- How the new techniques in machine learning, data mining and analytics can be leveraged in improving network and user experience?

Network slicing and mutli-domain E2E service management and orchestration

- > How to change the E2E service definition on the fly?
- How to deliver a network service offerings optimized for each and every use case, application and user?



